

## A GENERALIZATION OF SHELAH'S OMITTING TYPES THEOREM

By

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**Abstract.** This note gives generalizations of Shelah's omitting types theorem and Lopez–Escobar's Theorem.

### 1. Introduction

The omitting types theorem states that for a given countable set  $S$  of nonisolated types in  $T$ , there is a model of  $T$  omitting all the members of  $S$ , where  $T$  is a theory of a countable language. If  $L$  is uncountable, it is easy to construct an  $L$ -theory, that is, a counter example to the omitting types theorem. So, we are always interested in a theory with a countable language. There are many generalizations of the theorem. Among these, Shelah's omitting types theorem is of special interest.

**THEOREM (Shelah).** *Let  $T$  be a theory of a countable language  $L$ . Let  $R$  be a set of nonisolated complete types such that  $|R| < 2^\omega$ . Then there is a model  $M \models T$  omitting all the members of  $R$ .*

If we assume Martin's Axiom, we can omit  $< 2^\omega$  nonisolated types. Newelski studied the maximum cardinal  $\kappa$  such that we can omit  $< \kappa$  nonisolated types. It is known that there is a model of  $ZFC + \neg CH$  such that  $\kappa = \omega_1$  (see [4]). So, we cannot omit the assumption of the completeness of types in Shelah's omitting types theorem.

One of the main theorems in this paper is the following; it simultaneously generalizes the usual omitting types theorem and Shelah's omitting types theorem, and is proved in section 3.

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**THEOREM.** *Let  $T$  be a theory formulated in a countable language  $L$  and  $L_0$  a sublanguage of  $L$ . Let  $R$  be a set of nonisolated complete  $L_0$ -types such that  $|R| < 2^\omega$ . Let  $S$  be a countable set of nonisolated  $L$ -types. Then there is a model  $M \models T$  omitting all the members of  $R \cup S$ .*

In section 4, we apply the above theorem to another version of the omitting types theorem, the Lopes–Escobar theorem [3]. The Lopes–Escobar theorem is as follows.

**THEOREM (Lopez-Escobar).** *Let  $T$  be a theory formulated in a countable language  $L$  having a binary relation  $<$ . Let  $S$  be a countable set of  $L$ -types. Suppose that for any  $\alpha < \omega_1$ , there is a model  $M_\alpha$  of  $T$  omitting  $S$  and with the order type  $\alpha$ . Then there is a model  $N \models T$  omitting  $S$  and with the order type  $\mathbf{Q}$ .*

This theorem has already been generalized for uncountably many complete types by Tsuboi [8]. We generalize the theorem to our situation.

**THEOREM.** *Let  $T$  be a theory formulated in a countable language  $L$  and  $L_0$  a sublanguage of  $L$ , which have a binary relation  $<$ . Let  $R$  be a set of nonisolated complete  $L_0$ -types such that  $|R| < 2^\omega$ . Let  $S$  be a countable set of  $L$ -types. Suppose that for any  $\alpha < \omega_1$ , there is a model  $M_\alpha$  of  $T$  omitting all the members of  $R \cup S$  and with the order type  $\alpha$ . Then there is a model  $N \models T$  omitting all the members of  $R \cup S$  and with the order type  $\mathbf{Q}$ .*

The omitting types theorem is also studied in nonclassical logics,  $L_{\omega_1, \omega}$  [2],  $L(Q)$  [5], etc. Our generalization of the omitting types theorem implies that Shelah’s omitting types theorem holds in  $PC_\delta$ -classes. Precise definitions are given in section 5.

## 2. Preliminaries and Notations

Throughout,  $L$  is a countable language and  $T$  is a countable first-order theory formulated in  $L$ . ( $T$  may be incomplete.) We always work under  $T$ .  $L$ -formulas are denoted by  $\varphi, \psi, \theta, \chi, \dots$ . We fix a sublanguage  $L_0 \subset L$ .  $L_0$ -formulas are denoted by  $\zeta, \dots$ . We assume that  $\varphi, \psi, \dots, \zeta, \dots$  are satisfiable unless otherwise noted. Types are (possibly incomplete)  $L$ -types over the empty set. We say a type  $p(\bar{x})$  is a complete  $L_0$ -type if  $p$  consists of only  $L_0$ -formulas, and if for every  $\zeta(\bar{x}) \in L_0$ ,  $\zeta$  or  $\neg\zeta$  is in  $p$ .

EXAMPLE 1. Let  $L = \{R_i(x) \mid i < \omega\}$ . Consider an  $L$ -theory  $T$ , where for every finite subset  $F, G \subset \omega$  with  $F \cap G = \emptyset$ , there is an element  $x$  satisfying  $\bigwedge_{i \in F} R_i(x) \wedge \bigwedge_{j \in G} \neg R_j(x)$ . Then  $T$  is complete and not small. Moreover, there is no isolated complete type. Let  $R$  be a set of complete types with  $|R| < 2^\omega$ . Shelah's omitting types theorem says that there is a model  $M \models T$  omitting all the members of  $R$ . Take an infinite-coinfinite subset  $S \subset \omega$ . Set  $\Sigma(x) = \{R_i(x) \mid i \in S\}$ . Then

1.  $\Sigma$  is a nonprincipal type,
2.  $\Sigma$  has continuum many extensions to nonprincipal complete types.

So, it is not clear that there is a model  $N \models T$  omitting all the members of  $R \cup \{\Sigma\}$ .

### 3. A Proof of the Theorem

The main idea of the proof is simple; construct continuum many models such that each type in  $R \cup S$  is omitted by almost models. Then, there must be a model that omits all the members of  $R \cup S$  because  $|R \cup S| < 2^\omega$ . To prove the theorem, we make the following definitions.

DEFINITION 2. Let  $L_0 \subset L$  and  $\varphi_i(\bar{x}) \in L$ .

1. We say that two  $L$ -formulas  $\varphi_0(\bar{x})$  and  $\varphi_1(\bar{x})$  are  $L_0$ -separable in  $\bar{x}' \subset \bar{x}$  if there are  $L_0$ -formulas  $\xi_0(\bar{x}')$  and  $\xi_1(\bar{x}')$  such that  $T \models \varphi_k(\bar{x}) \rightarrow \xi_k(\bar{x}')$  ( $k = 0, 1$ ), and  $\xi_0$  and  $\xi_1$  are incompatible in  $T$ .
2. We say  $\varphi_0(\bar{x})$  and  $\varphi_1(\bar{x})$  are essentially  $L_0$ -separable in  $\bar{x}' \subset \bar{x}$  if there are  $L$ -formulas  $\varphi'_k(\bar{x})$  ( $k = 0, 1$ ) with  $T \models \varphi'_k(\bar{x}) \rightarrow \varphi_k(\bar{x})$  ( $k = 0, 1$ ) such that  $\varphi'_0$  and  $\varphi'_1$  are  $L_0$ -separable in  $\bar{x}'$ .
3. Let  $\Phi = \varphi_0(\bar{x}), \dots, \varphi_n(\bar{x})$  be a sequence of  $L$ -formulas. We say that  $\Phi$  is maximally  $L_0$ -separated if for each  $i \neq j$  and each subsequence  $\bar{x}' \subset \bar{x}$ , whenever  $\varphi'_i(\bar{x})$  and  $\varphi'_j(\bar{x})$  are essentially  $L_0$ -separable in  $\bar{x}'$  then they are  $L_0$ -separable in  $\bar{x}'$ .

A maximally  $L_0$ -separated sequence  $\Phi' = \varphi'_0(\bar{x}), \dots, \varphi'_n(\bar{x})$  will be called a maximal  $L_0$ -separation of  $\Phi$  if  $T \models \varphi'_i(\bar{x}) \rightarrow \varphi_i(\bar{x})$  ( $i = 0, \dots, n$ ).

LEMMA 3. Let  $\Phi = \varphi_0(\bar{x}), \dots, \varphi_n(\bar{x})$  be  $L$ -formulas. Then there are  $L$ -formulas  $\varphi'_i(\bar{x})$  ( $i \leq n$ ) such that  $\Phi' = \varphi'_0(\bar{x}), \dots, \varphi'_n(\bar{x})$  is a maximal  $L_0$ -separation of  $\Phi$ .

PROOF. Let  $\bar{y} \subset \bar{x}$  and suppose that  $\varphi_i(\bar{y})$  and  $\varphi_j(\bar{y})$  are essentially  $L_0$ -separable in  $\bar{y}$ . Choose an  $L$ -formula  $\varphi'_i(\bar{x})$  and an  $L$ -formula  $\varphi'_j(\bar{x})$  witnessing the

essential  $L_0$ -separability. Then we replace  $\varphi_i(\bar{x})$  and  $\varphi_j(\bar{x})$  by  $\varphi'_i(\bar{x})$  and  $\varphi'_j(\bar{x})$ , respectively. We repeat this process (finitely many times) and finally we get a desired maximal  $L_0$ -separation.

DEFINITION 4. Let  $\psi(x_1, \dots, x_n)$  be an  $L$ -formula and  $s(\bar{y})$  an  $L$ -type. We say  $\psi(x_1, \dots, x_n)$  totally omits  $s(\bar{y})$  if whenever  $M \models T$  and  $a_1, \dots, a_n \in M$  satisfy  $\psi(\bar{x})$  then no tuple from  $\{a_1, \dots, a_n\}$  realizes  $s(\bar{y})$ . Let  $\Sigma$  be a finite set of formulas. We simply say that  $\Sigma$  totally omits  $s$  if  $\bigwedge \Sigma$  totally omits  $s$ .

- REMARK 5. • Let  $s(\bar{x})$  be a nonisolated type. Then for every  $L$ -formula  $\varphi(\bar{x})$  there is an  $L$ -formula  $\varphi'(\bar{x})$  with  $T \models \varphi'(\bar{x}) \rightarrow \varphi(\bar{x})$  such that  $\varphi'$  and  $s$  are inconsistent.
- It is easy to check that for every  $L$ -formula  $\varphi(\bar{x})$  and nonisolated type  $s(\bar{y})$ , there is an  $L$ -formula  $\psi(\bar{x})$  with  $T \models \psi \rightarrow \phi$  such that  $\psi$  totally omits  $s$ .

Next lemma is easy but important for our proof of the theorem.

LEMMA 6. Let  $\varphi_0(\bar{x})$  and  $\varphi_1(\bar{x})$  be  $L$ -formulas such that they are not essentially  $L_0$ -separable in  $\bar{x}' \subset \bar{x}$ . Then  $\varphi_0$  and  $\varphi_1$  isolate the same complete  $L_0$ -type  $p(\bar{x}')$ .

PROOF. Suppose otherwise. Then it is easy to find an  $L_0$ -formula  $\zeta(\bar{x}')$  such that both  $\varphi_0 \wedge \zeta$  and  $\varphi_1 \wedge \neg \zeta$  are satisfiable. Two  $L$ -formulas  $\varphi_0 \wedge \zeta$  and  $\varphi_1 \wedge \neg \zeta$  are  $L_0$ -separable in  $\bar{x}'$ . Since  $T \models \varphi_0 \wedge \zeta \rightarrow \varphi_0$  and  $T \models \varphi_1 \wedge \neg \zeta \rightarrow \varphi_1$ , this means that  $\varphi_0$  and  $\varphi_1$  are essentially  $L_0$ -separable. A contradiction.

THEOREM 7. Let  $R$  be a set of nonisolated complete  $L_0$ -types such that  $|R| < 2^\omega$ . Let  $S$  be a countable set of nonisolated  $L$ -types. Then there is a model  $M \models T$  omitting all the members of  $R \cup S$ .

PROOF. Suppose  $Z = \{z_i \mid i < \omega\}$  is a fixed countable set of new variables. We denote a sequence  $z_0, z_1, \dots, z_{i-1}$  by  $\bar{z}_i$ . Enumerate  $S$  as  $S = \{s_i(\bar{x}_i) : i \in \omega\}$ . We may assume that for each  $s_n(\bar{x}_n)$ ,  $|\bar{x}_n| \leq n$ . Let  $\{\theta_i(\bar{z}_i, z_i)\}$  be an enumeration of the  $L$ -formulas having the form  $\exists x \varphi(\bar{z}_i, x) \rightarrow \varphi(\bar{z}_i, z_i)$ .

By induction, we construct a binary tree  $\{\Sigma_\eta(\bar{z}_{len(\eta)}) \mid \eta \in 2^{<\omega}\}$  of finite sets of  $L$ -formulas with the following properties: For every  $n \in \omega$  and every  $\eta \in 2^n$ ,

1. If  $m < n$  then  $\Sigma_{\eta|m} \subset \Sigma_{\eta|n}$ ;
2.  $\{\bigwedge \Sigma_\sigma(\bar{z}_n)\}_{\sigma \in 2^n}$  is maximally separated;

3.  $\Sigma_\eta$  is consistent;
4.  $\Sigma_\eta$  contains  $\theta_n$ ;
5.  $\Sigma_\eta$  totally omits each of  $s_i$  ( $i \leq n$ ).

Let  $\Sigma_{\langle \cdot \rangle} = \emptyset$  and suppose  $\Sigma_\sigma(\bar{z}_n)$  has been defined for every  $\sigma \in 2^n$ . Take two copies of  $\Sigma_\sigma(\bar{z}_n)$  and set

$$\Sigma_\sigma^{0,k}(\bar{z}_n) = \Sigma_\sigma(\bar{z}_n) \quad (k = 0, 1).$$

Then, by Lemma 3, there is a set  $\{\psi_{\sigma,k}(\bar{z}_n)\}_{\sigma \in 2^n, k=0,1}$  which is a maximal  $L_0$ -separation of  $\{\bigwedge \Sigma_\sigma^{0,k}(\bar{z}_n)\}_{\sigma \in 2^n, k=0,1}$ . Set

$$\Sigma_\sigma^{1,k}(\bar{z}_n) = \Sigma_\sigma^{0,k}(\bar{z}_n) \cup \{\psi_{\sigma,k}(\bar{z}_n)\}.$$

Next, for each  $\sigma \in 2^n$ , take an  $L$ -formula  $\chi_{\sigma,k}(\bar{z}_n) \models \Sigma_\sigma^{1,k}(\bar{z}_n)$  such that  $\chi_{\sigma,k}$  totally omits  $s_i(\bar{x}_i)$  for every  $i \leq n$ . (Such formula exists by Remark 5.) Set

$$\Sigma_\sigma^{2,k}(\bar{z}_n) = \Sigma_\sigma^{1,k}(\bar{z}_n) \cup \{\chi_{\sigma,k}(\bar{z}_n)\}.$$

Finally set  $\Sigma_{\sigma^{\cdot k}} = \Sigma_\sigma^{2,k}(\bar{z}_n) \cup \{\theta_n(\bar{z}_n, z_n)\}$ . It is easy to check that  $\{\Sigma_\eta(\bar{z}_{n+1})\}_{\eta \in 2^{n+1}}$  satisfies the required conditions 1–5 (with  $n$  replaced by  $n+1$ ). So we have succeeded to construct all  $\Sigma_\eta$ 's. Now, for a path  $\eta \in 2^\omega$ , we define  $\Sigma_\eta(Z)$  by  $\Sigma_\eta = \bigcup_{n \in \omega} \Sigma_{\eta|n}$ . Recall that  $\theta_n$  has the form  $\exists x \varphi(\bar{z}_n, x) \rightarrow \varphi(\bar{z}_n, z_n)$ . So, by the condition 4, every  $M_\eta$  realizing  $\Sigma_\eta(Z)$  is a model of  $T$ . By the condition 5,  $M_\eta$  omits all types in  $S$ .

CLAIM A. *For each  $p \in R$ ,  $\{\eta \in 2^\omega \mid M_\eta \models \exists \bar{x} p(\bar{x})\}$  is countable.*

We fix  $p(\bar{x}) \in R$  and  $\bar{z} \subset Z$  with  $|\bar{x}| = |\bar{z}|$ . Suppose  $\Sigma_\eta(Z) \cup p(\bar{z})$  is consistent. Take any  $\eta' \neq \eta$ . If  $\Sigma_{\eta'}(Z) \cup p(\bar{z})$  is also consistent, then  $\Sigma_{\eta|n}$  and  $\Sigma_{\eta'|n}$  are not essentially  $L_0$ -separable in  $\bar{z}$ , where  $n$  is chosen so that  $\bar{z} \subset \bar{z}_n$ . Hence  $p$  must be isolated by a  $L$ -formula, by Lemma 6. But  $R$  is a set of nonisolated types, a contradiction. So, for each  $p \in R$  and  $\bar{z} \subset Z$ ,  $\{\eta \in 2^\omega \mid \Sigma_\eta(Z) \cup p(\bar{z}) \text{ consistent}\}$  has at most one element. This proves the claim, since there are only countably many possible choices of  $\bar{z} \subset Z$ . (End of Proof of Claim)

Finally, by the claim above and the assumption that  $|R| < 2^\omega$ , we can find a path  $\eta \in 2^\omega$  such that  $M_\eta$  omits  $R$ .

COROLLARY 8. *Suppose  $\alpha < 2^\omega$ . Let  $T_0$  be a complete  $L$ -theory and  $p, q_i \in S(T_0)$  ( $i < \alpha$ ). If for every  $i < \alpha$  there is a model  $M_i$  such that  $M_i$  omits  $q_i$  and  $M_i$  realizes  $p$ , then there is a model  $N$  such that  $N$  omits all  $q_i$ 's but  $N$  realizes  $p$ .*

#### 4. Another Version of Omitting Types Theorem with Uncountably Many Types

Recall that  $L$  is a countable language and  $L_0$  a sublanguage of  $L$ . In this section we show the following,

**THEOREM 9.** *Let  $T$  be a (possibly incomplete)  $L$ -theory. Let  $R$  be a set of complete  $L_0$ -types with  $|R| < 2^\omega$  and  $S$  a countable set of  $L$ -types. Fix an  $L$ -formula  $\chi(x, y)$ . Suppose that for any  $\alpha < \omega_1$ , there is a model  $M_\alpha$  of  $T$  containing a set  $A_\alpha \subset M_\alpha$  such that*

- $A_\alpha = \{a_i^\alpha \mid i \leq \alpha\}$ ,
- $M_\alpha \models \chi(a_i^\alpha, a_j^\alpha)$  if and only if  $i < j$ ,
- $M_\alpha$  omits all the members of  $R \cup S$ .

Then there is a model  $N \models T$  with a subset  $A \subset N$  such that

- $A = \{a_q \mid q \in \mathbf{Q}\}$ ,
- $N \models \chi(a_q, a_{q'})$  if and only if  $q < q'$ ,
- $N$  omits all the members of  $R \cup S$ .

In the rest of this section, we denote  $\chi(x, y)$  by  $x < y$ . For a tuple  $\bar{a}$ , the  $i$ th element of  $\bar{a}$  is denoted by  $(\bar{a})_i$ . We also denote the  $\beta$ th element  $a_{i+\beta}^\gamma$  from  $a_i^\gamma$  in  $A_\gamma$  by  $a_i^\gamma + \beta$ .

Note that if, with new constants  $c_q$  ( $q \in \mathbf{Q}$ ),  $T \cup \{c_q < c_{q'} \mid q < q', q, q' \in \mathbf{Q}\}$  isolates no type in  $R \cup S$  then the theorem is clear by theorem 7. But, in general,  $T \cup \{c_q < c_{q'}\}_{q, q'}$  may isolate some types. (Notice that  $T \cup \{c_q < c_{q'}\}_{q, q'}$  may not be complete.) So we need find a theory  $T' \supset T$  that isolates no type in  $R \cup S$ . To construct  $T'$ , we need some definitions. The following definitions are taken from the proof of Lopez–Escobar’s theorem in [2].

- DEFINITION 10.**
1. An  $m$ -sequence is a sequence of tuples of length  $m$ .
  2. We say an ascending tuple  $\bar{b} \in A_\gamma$  of length  $m + 1$  is a  $k$ -extension ( $k \leq m$ ) of an ascending tuple  $\bar{a} \in A_\gamma$  of length  $m$  if  $(\bar{b})_1 = (\bar{a})_1, \dots, (\bar{b})_k = (\bar{a})_k, (\bar{b})_{k+2} = (\bar{a})_{k+1}, \dots, (\bar{b})_{m+1} = (\bar{a})_m$ .
  3. Let  $\Gamma$  be a subset of  $\omega_1$ . We say that the  $m$ -sequence  $\{\bar{a}^\gamma \mid \gamma \in \Gamma\}$  is an unbounded  $m$ -sequence if
    - $\Gamma$  is unbounded in  $\omega_1$ ,
    - $\bar{a}^\gamma$  is an ascending tuple of length  $m$  of elements of  $A_\gamma$ ,

- for any  $\beta \in \omega_1$  there is a  $\gamma \in \Gamma$  such that  $a_\beta^\gamma < (\bar{a}^\gamma)_1$ ,  $(\bar{a}^\gamma)_1 + \beta < (\bar{a}^\gamma)_2$ ,  $(\bar{a}^\gamma)_2 + \beta < (\bar{a}^\gamma)_3, \dots, (\bar{a}^\gamma)_m + \beta < a_\beta^\gamma$ .
- 4. Let  $\Gamma$  be a subset of  $\omega_1$ . Let  $X = \{\bar{a}^\gamma \mid \gamma \in \Gamma\}$  be an unbounded  $m$ -sequence and  $Y = \{\bar{b}^\gamma \mid \gamma \in \Gamma\}$  an unbounded  $(m+1)$ -sequence. We say  $Y$  is a  $k$ -extension of  $X$  ( $0 \leq k \leq m$ ) if for all  $\gamma \in \Gamma$ ,  $\bar{b}^\gamma$  is a  $k$ -extension of  $\bar{a}^\gamma$ .
- 5. We consider the unbounded 0-sequence, the empty sequence. Every unbounded 1-sequence is a 0-extension of the unbounded 0-sequence.

LEMMA 11. 1. *There is an unbounded 1-sequence.*

2. *Let  $X$  be an unbounded  $m$ -sequence and  $k \leq m$ . Then there are an unbounded  $(m+1)$ -sequence  $Y$  and an unbounded  $m$ -sequence  $X'$  such that  $X'$  is an unbounded  $m$ -sequence,  $X' \subset X$ , and  $Y$  is a  $k$ -extension of  $X'$ . This condition will be denoted as  $X \triangleleft_k Y$ .*

PROOF. We show the second with  $m = 1$  and  $k = 0$ , and the other cases are similar. Let  $X = \{a^\gamma\}_{\gamma \in \Gamma}$  be an unbounded 1-sequence. Then for any  $\beta + \beta + \beta \in \omega_1$  there is a  $\gamma \in \Gamma$  such that  $a_{\beta,3}^\gamma < a^\gamma$  (Recall  $a_{\beta,3}^\gamma$  is the  $\beta \cdot 3$ -th element of  $A_\gamma$ ). So we have a 0-extension  $\bar{b}^\gamma = a_\beta^\gamma + 1$ ,  $a^\gamma$  of  $a^\gamma$ . Collect such 0-extension  $\bar{b}^\gamma$  of  $a^\gamma$  for every  $\beta \in \omega$ , then it is a required 2-sequence.

Take a set  $C = \{c_q \mid q \in \mathbf{Q}\}$  of new constant symbols. To prove the theorem, it is enough to show that there is an  $L$ -theory  $T' \supset T \cup \{c_q < c_{q'} \mid q < q' \text{ and } q, q' \in \mathbf{Q}\}$  such that all the members of  $R$  and  $S$  are nonisolated types in  $T'$ , by theorem 7. We fix an enumeration  $\{c_{q_n} \mid n < \omega\}$  of  $C$ . Let  $\bar{c}_n$  be the sequence consisting  $c_{q_0}, c_{q_1}, \dots, c_{q_{n-1}}$  with the order of  $\mathbf{Q}$  (e.g. if  $q_0, q_1, q_2 = 0.5, -1, 0$  then  $\bar{c}_3$  is the sequence  $c_{-1}, c_0, c_{0.5}$ ). Most ideas of the following definitions are from [8]. We adapt it to our situation.

DEFINITION 12. Let  $X = \{\bar{a}^\gamma \mid \gamma \in \Gamma\}$  be an unbounded  $m$ -sequence with  $\Gamma \subset \omega_1$  and  $\varphi(\bar{x}, \bar{c})$  an  $L(\bar{c})$ -formula.

1. We say  $X$  is  $\varphi(\bar{x}, \bar{c})$ -uniform if for every  $L_0$ -formula  $\xi(\bar{x})$  and  $\gamma, \gamma' \in \Gamma$ ,  $M_\gamma \models \exists \bar{x}(\varphi(\bar{x}, \bar{a}^\gamma) \wedge \xi(\bar{x}))$  if and only if  $M_{\gamma'} \models \exists \bar{x}(\varphi(\bar{x}, \bar{a}^{\gamma'}) \wedge \xi(\bar{x}))$
2. We say  $X$  is essentially  $\varphi(\bar{x}, \bar{c})$ -uniform if there is an unbounded subset  $\Gamma' \subset \Gamma$  such that  $X'$  is  $\varphi(\bar{x}, \bar{c})$ -uniform where  $X' = \{\bar{a}^\gamma \in X \mid \gamma \in \Gamma'\}$ .

LEMMA 13. *Let  $X = \{\bar{a}^\gamma \mid \gamma \in \Gamma\}$  be an unbounded  $m$ -sequence with  $\Gamma \subset \omega_1$  and  $\varphi(\bar{x}, \bar{c})$  an  $L(\bar{c})$ -formula. If  $X$  is not essentially  $\varphi(\bar{x}, \bar{c})$ -uniform then there is an*

$L_0$ -formula  $\xi(\bar{x})$  such that  $X_{\varphi \wedge \xi}$  and  $X_{\varphi \wedge \neg \xi}$  are unbounded  $m$ -sequences where  $X_{\theta(\bar{x}, \bar{c})} := \{\bar{a}' \mid M_\gamma \models \exists \bar{x} \theta(\bar{x}, \bar{a}')\}$ .

PROOF. Suppose that  $X_{\varphi \wedge \xi}$  or  $X_{\varphi \wedge \neg \xi}$  is bounded for every  $L_0$ -formula  $\xi(\bar{x})$ . Notice that if  $X_{\varphi \wedge \xi}$  is bounded then  $X_{\varphi \wedge \xi}$  is countable. So, the union  $Y$  of all bounded  $X_{\varphi \wedge \xi}$ 's is also countable, because  $L_0$  is countable. Set  $X' = X \setminus Y$ . Then  $X'$  is an unbounded  $m$ -sequence, and  $X'$  is  $\varphi(\bar{x}, \bar{c})$ -uniform by the definition of  $X'$ . This means that  $X$  is essentially  $\varphi$ -uniform.

LEMMA 14. *Let  $Y$  and  $Y'$  be unbounded  $m$ -sequences. Suppose they are not essentially  $\varphi_i(\bar{x})$ -uniform for  $i \leq n$ . Then there is an  $L_0$ -formula  $\xi_i(\bar{x})$ ,  $X \subset Y$  and  $X' \subset Y'$  such that  $X_{\varphi_i \wedge \xi_i} = X$ ,  $X'_{\varphi_i \wedge \neg \xi_i} = X'$  and  $X, X'$  are unbounded  $m$ -sequences, for each  $i \leq n$ .*

PROOF. We show by induction on  $n$ . The case  $n = 0$  is trivial. Let  $n = k + 1$ . By induction hypothesis, we have  $Z \subset Y$  and  $Z' \subset Y'$  such that  $Z_{\varphi_i \wedge \xi_i} = Z$ ,  $Z'_{\varphi_i \wedge \neg \xi_i} = Z'$  for each  $i \leq k$ . Let  $\xi^0$  be an  $L_0$ -formula dividing  $Z$  into two uncountable sets  $Z_{\varphi_{k+1} \wedge \xi^0}$ ,  $Z_{\varphi_{k+1} \wedge \neg \xi^0}$ . By shrinking  $Z'$ , if necessary, we may assume  $Z'_{\varphi_{k+1} \wedge \neg \xi^0} = Z'$ . Then, let  $\xi^1$  be an  $L_0$ -formula dividing  $Z'$  into two uncountable sets  $Z'_{\varphi_{k+1} \wedge \xi^1}$ ,  $Z'_{\varphi_{k+1} \wedge \neg \xi^1}$ . Either  $Z_{\varphi_{k+1} \wedge \xi^0 \wedge \xi^1}$  or  $Z_{\varphi_{k+1} \wedge \xi^0 \wedge \neg \xi^1}$  is uncountable, we can take  $\xi^0 \wedge \xi^1$  or  $\xi^0 \wedge \neg \xi^1$  as  $\xi_{k+1}$ . Then put  $X = Z_{\varphi_{k+1} \wedge \xi_{k+1}}$  and  $X' = Z'_{\varphi_{k+1} \wedge \neg \xi_{k+1}}$ .

Let  $\{\varphi_n(\bar{x}, \bar{c}_n) \mid n < \omega\}$  be an enumeration of all  $L(C)$ -formulas. Also enumerate  $S$  as  $S = \{s_n(\bar{x}_n) \mid n < \omega\}$ . We can assume that for every tuple  $(\varphi, s) \in L(C) \times S$ , there is  $n$  such that  $(\varphi, s) = (\varphi_n, s_n)$ . So, each member of  $L(C)$ ,  $S$  appears infinitely many times in the enumerations. By induction, we construct a binary tree  $\{T^\sigma(\bar{c}_{len(\sigma)}) \mid \sigma \in 2^{<\omega}\}$  of sets of  $L(C)$ -formulas and unbounded  $len(\sigma)$ -sequence  $X^\sigma = \{\bar{a}_\gamma \mid \gamma \in \Gamma_\sigma\}$  with the following properties: For every  $\sigma, \sigma' \in 2^{<\omega}$  and  $n \leq len(\sigma)$ ,

1.  $T^\sigma(\bar{c}_{len(\sigma)}) \cup \{c_q < c_{q'} \mid c_q, c_{q'} \in \bar{c}_{len(\sigma)} \text{ and } q < q'\}$  is consistent,
2.  $\sigma \subset \sigma'$  then  $T^\sigma \subset T^{\sigma'}$ ,
3.  $M_\gamma, \bar{a}_\gamma \models T^\sigma(\bar{c}_{len(\sigma)})$  for uncountably many  $\gamma \in \Gamma_\sigma$ ,
4.  $T^\sigma$  contains  $\exists \bar{x} \varphi_{len(\sigma)}$  or  $\neg \exists \bar{x} \varphi_{len(\sigma)}$ ,
5.  $T^\sigma$  contains  $\exists \bar{x} (\varphi_{len(\sigma)}(\bar{x}, \bar{c}_{len(\sigma)}) \wedge \neg \psi(\bar{x}))$  for some  $\psi \in s_{len(\sigma)}$ ,
6. if  $X^\sigma$  is essentially  $\varphi_n$ -uniform then it is  $\varphi_n$ -uniform,
7. if  $X^\sigma$  is not essentially  $\varphi_n$ -uniform then  $\exists \bar{x} (\xi(\bar{x}) \wedge \varphi_n(\bar{x}, \bar{c}_{len(\sigma)})) \in T^{\sigma^0}$  and  $\exists \bar{x} (\neg \xi(\bar{x}) \wedge \varphi_n(\bar{x}, \bar{c}_{len(\sigma)})) \in T^{\sigma^1}$  for some  $\xi \in L_0$ .



Let  $T^{\langle \rangle} = \emptyset$ ,  $X^{\langle \rangle}$  the unbounded 0-sequence and suppose  $T^\sigma(\bar{c}_n)$  and  $X^\sigma$  are defined for every  $\sigma \in 2^n$ . Suppose  $\bar{c}_{n+1}$  is a  $k$ -extension of  $\bar{c}_n$ . Take an unbounded  $(n+1)$ -sequence  $Y^\sigma \triangleright_k X^\sigma$  (Lemma 11). If  $Y_{\varphi_{n+1}}^\sigma$  is uncountable, set

$$Y^{\sigma,0} = Y_{\varphi_{n+1}}^\sigma$$

and

$$T^{\sigma,0} = T^\sigma \cup \{\exists \bar{x} \varphi_{n+1}(\bar{x}, \bar{c}_{n+1})\}$$

otherwise

$$Y^{\sigma,0} = Y_{\neg \varphi_{n+1}}^\sigma,$$

$$T^{\sigma,0} = T^\sigma \cup \{\neg \exists \bar{x} \varphi_{n+1}(\bar{x}, \bar{c}_{n+1})\}.$$

Recall that  $s_{n+1}(\bar{x}_{n+1})$  is countable and  $M_\alpha$  omits  $s_{n+1}(\bar{x}_{n+1})$  for every  $\alpha$ . Hence, we can find  $\psi(\bar{x}_{n+1}) \in s_{n+1}$  such that  $Y_{\varphi_{n+1} \wedge \neg \psi}^{\sigma,0}$  is uncountable. Set

$$Y^{\sigma,1} = Y_{\varphi_{n+1} \wedge \neg \psi}^{\sigma,0},$$

$$T^{\sigma,1} = T^{\sigma,0} \cup \{\exists \bar{x}(\varphi_{n+1}(\bar{x}, \bar{c}_{n+1}) \wedge \psi(\bar{x}))\}.$$

Then, if  $Y^{\sigma,1}$  is essentially  $\varphi_{n+1}$ -uniform, by shrinking it, we may assume  $Y^{\sigma,1}$  is  $\varphi_{n+1}$ -uniform.

Finally, we consider the  $\varphi_j$ -uniformity of  $Y^{\sigma,1}$  ( $j \leq N+1$ ). If  $Y^{\sigma,1}$  is  $\varphi_j$ -uniform for every  $j \leq n+1$  then set  $X^{\sigma^0} = X^{\sigma^1} = Y^{\sigma,1}$  and  $T^{\sigma^0} = T^{\sigma^1} = T^{\sigma,1}$ . Otherwise, assume  $Y^{\sigma,1}$  is not essentially  $\varphi_j$ -uniform for some  $j$ . Then, take an unbounded  $(n+1)$ -sequence  $X^{\sigma^i} \subset Y^{\sigma,1}$  ( $i = 0, 1$ ) such that for all  $j \leq n$ , if  $Y^{\sigma,1}$  is not essentially  $\varphi_j$ -uniform, then  $X_{\varphi_j \wedge \xi_j}^{\sigma^0} = X^{\sigma^0}$  and  $X_{\varphi \wedge \neg \xi_j}^{\sigma^1} = X^{\sigma^1}$  for some  $\xi_j(\bar{x}) \in L_0$  (See lemma 14). We set

$$T^{\sigma^0} = T^{\sigma,1} \cup \{\exists \bar{x}(\varphi_j(\bar{x}, \bar{c}_j) \wedge \xi_j(\bar{x}))\}_j,$$

$$T^{\sigma^1} = T^{\sigma,1} \cup \{\exists \bar{x}(\varphi_j(\bar{x}, \bar{c}_j) \wedge \neg \xi_j(\bar{x}))\}_j.$$

It is easy to check that they satisfy the required conditions. At the end of this inductive construction, we have  $2^\omega$  complete  $L(C)$ -theories  $T^\eta$  ( $\eta \in 2^\omega$ ). By condition 5. and the way of enumerations of  $L(C)$  and  $S$ , every member of  $S$  is not isolated in  $T^\eta$ .

CLAIM A. *The set  $\{\eta \in 2^\omega \mid p \text{ is isolated in } T^\eta\}$  is countable for every  $p \in R$ .*

Suppose  $p(\bar{x})$  is isolated by an  $L(C)$ -formula  $\varphi_n(\bar{x}, \bar{c}_n)$  in  $T^\eta$  and  $T^{\eta'}$ . If  $X^{\eta^n} = \{\bar{a}_\gamma \mid \gamma \in \Gamma_{\eta^n}\}$  is  $\varphi_n$ -uniform then  $M_\gamma$  ( $\gamma \in \Gamma_{\eta^n}$ ) realizes  $p(\bar{x})$ . So,  $X^{\eta^n}$  is

not essentially  $\varphi_n$ -uniform. We can assume that  $\eta|n \neq \eta'|n$ , because, by condition 7., if  $\eta|n = \eta'|n$ ,  $\varphi_n$  cannot isolate same complete  $L_0$ -type in  $T^\eta$ ,  $T^{\eta'}$ . Therefore,  $\{\eta \in 2^\omega \mid p \text{ is isolated in } T^\eta\}$  is countable. (End of proof of claim)

Hence, there is an  $\eta \in 2^\omega$  such that every member of  $R$  is nonisolated in  $T^\eta$ . By theorem 7, we have a required model. (End of proof of theorem 9)

### 5. Omitting Types Theorem with Nonelementary Classes

In this section, we look at the definitions of some nonelementary classes. Then we have Shelah's omitting types theorem for such classes.

DEFINITION 15. Let  $\mathcal{K}$  be a class of  $L$ -structures. We say  $\mathcal{K}$  is an  $EC(\aleph_0, \aleph_0)$ -class if

- $L$  is countable,
- there is a countable set  $S$  of types and an  $L$ -theory  $T$  such that  $M \in \mathcal{K}$  if and only if  $M \models T$ , and  $M$  omits all the members of  $S$ .

$\mathcal{K}$  is denoted by  $EC(T, S)$ .

More general definitions and properties of  $EC(\kappa, \lambda)$  can be found in [1]. It is well known that every  $EC(\aleph_0, \aleph_0)$ -class can be translated to a class defined by an  $L_{\omega_1, \omega}$ -sentence, and vice versa (see [7], for example). Next, we introduce a  $PC_\delta$ -class. This is defined by Keisler in [2] with  $L_{\omega_1, \omega}$ . The following definition of a  $PC_\delta$ -class is given without  $L_{\omega_1, \omega}$ . Note that Shelah and Baldwin use other notations, e.g.,  $PC(\aleph_0, \aleph_0)$ ,  $PC\Gamma(\aleph_0, \aleph_0)$  (see [1]).

DEFINITION 16. Let  $\mathcal{K}$  be a class of  $L$ -structures. We say that  $\mathcal{K}$  is a  $PC_\delta$ -class if there is a countable language  $L' \supset L$  and a class of  $L'$ -structures  $\mathcal{K}'$  such that

- $\mathcal{K}'$  is an  $EC(\aleph_0, \aleph_0)$ -class, and
- $M' \upharpoonright L \in \mathcal{K}$  if and only if  $M' \in \mathcal{K}'$  for every  $L'$ -structure  $M'$ .

To generalize the omitting types theorem for a  $PC_\delta$ -class, we need definitions of types and isolated types.

DEFINITION 17. Let  $\mathcal{K}$  be a class of  $L$ -structures. A type  $\Sigma(\bar{x})$  in  $\mathcal{K}$  is a set of  $L$ -formulas with free variables  $\bar{x}$  such that there is a structure  $M \in \mathcal{K}$  having a realization of  $\Sigma(\bar{x})$ .

DEFINITION 18. Let  $\mathcal{K}$  be either  $EC(T, S)$  or the  $PC_\delta$ -class obtained from  $EC(T, S)$  by restricting the language. An isolated type in  $\mathcal{K}$  is a type in  $\mathcal{K}$  which is isolated in  $T$  in the usual sense.

The above definitions and theorem 7 immediately give Shelah's omitting types theorem for  $PC_\delta$ -classes.

THEOREM 19. *Let  $\mathcal{K}$  be a  $PC_\delta$ -class. Let  $R$  be a set of nonisolated complete types in  $\mathcal{K}$  such that  $|R| < 2^\omega$ . Then there is a model  $M \models T$  omitting all the members of  $R$ .*

We also have Lopez-Escobar's theorem (with uncountably many types) for  $PC_\delta$ -classes.

THEOREM 20. *Let  $\mathcal{K}$  be a  $PC_\delta$ -class with a countable language  $L$  having a binary relation  $<$ . Let  $R$  be a set of complete  $L$ -types such that  $|R| < 2^\omega$ . Suppose that for any  $\alpha < \omega_1$ , there is a model  $M_\alpha \in \mathcal{K}$  omitting  $R$  and with the order type  $\alpha$ . Then there is a model  $N \in \mathcal{K}$  omitting  $R$  and with the order type  $\mathbf{Q}$ .*

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