

type theory の回避
 集合論の paradox

$$X \in X \\ X \in \mathbb{N}$$

Russell & Whitehead

Principia Mathematica

⇓ 20世紀の数学
 1930年代

Church
 formal

homotopy type theory \Rightarrow dependent
 type theory

simple type theory

Categorical Logic and Type theory

B. Jacobs

$\{\mathbb{N}, \mathbb{B}, \dots\}$
 自然数 2-16進数値

$$+ : \mathbb{N}, \mathbb{N} \rightarrow \mathbb{N}$$

$$\text{succ} : \mathbb{N} \rightarrow \mathbb{N} \quad n \mapsto n+1$$

$$\wedge : \mathbb{B}, \mathbb{B} \rightarrow \mathbb{B} \quad \text{and}$$

$$= : \mathbb{N}, \mathbb{N} \rightarrow \mathbb{B}$$

signature

$$\Sigma = (T, F)$$

T: a set of Basic type

F: 関数の集合

calculus 変数 $\{x_1, x_2, \dots\}$ 可算個

context

context P

$$P(x_1 : A_1, x_2 : A_2, \dots, x_n : A_n)$$

$$\Delta(x_1 : B_1, x_2 : B_2, \dots, x_m : B_m)$$

concatenation

$$P, \Delta = (x_1 : A_1, \dots, x_n : A_n, x_{n+1} : B_1, \dots, x_{n+m} : B_m)$$

term

$$\Gamma \vdash M : A \quad \langle k \rangle \in J$$

$$n : \mathbb{N}, m : \mathbb{N} \vdash \text{plus} \text{ (times } (m, n), m) : \mathbb{N}$$

規則

$$\frac{}{\lambda : A \vdash \lambda : A} \text{ (identity)}$$

$$\frac{\Gamma \vdash M_1 : A_1, \dots, \Gamma \vdash M_n : A_n \quad F : A_1, \dots, A_n \rightarrow A_{n+1}}{\Gamma \vdash F(M_1, \dots, M_n) : A_{n+1}} \text{ (function symbol)}$$

$$\frac{\Gamma \vdash M_1 : A_1, \dots, \Gamma \vdash M_n : A_n \quad \lambda : A_1, \dots, A_n \vdash M : B}{\Gamma \vdash M_1, \dots, M_n, \lambda : A_1, \dots, A_n \vdash M : B} \text{ (weakening)}$$

$$\frac{\Gamma \vdash \lambda : A, \lambda : A \vdash M : B}{\Gamma \vdash \lambda : A \vdash M[\lambda/\lambda] : B} \text{ (contraction)}$$

$$\frac{\Gamma, \lambda_i : A_i, \lambda_{i+1} : A_{i+1}, \Delta \vdash M : B}{\Gamma, \lambda_i : A_{i+1}, \lambda_{i+1} : A_i \Delta \vdash M[\lambda_i/\lambda_{i+1}, \lambda_{i+1}/\lambda_i] : B} \text{ (exchange)}$$

$$\text{plus} : \mathbb{N}, \mathbb{N} \rightarrow \mathbb{N}$$

$$\text{if} : \mathbb{B}, \mathbb{N}, \mathbb{N} \rightarrow \mathbb{N}$$

$$\lambda_1 : \mathbb{B}, \lambda_2 : \mathbb{N} \vdash \text{if} (\lambda_1, \lambda_2, \text{plus} (\lambda_2, \lambda_2)) : \mathbb{N}$$

derivable

証)

$$\frac{\frac{K_1 : B \vdash K_1 : B}{K_1 : B, K_2 : B \vdash K_1 : B}}{\frac{K_1 : \mathbb{N} \vdash K_1 : \mathbb{N} \quad K_1 : \mathbb{N}, K_2 : \mathbb{N} \vdash K_1 : \mathbb{N}}{K_1 : \mathbb{N}, K_2 : \mathbb{N} \vdash K_2 : \mathbb{N}} \quad \frac{K_1 : \mathbb{N} \vdash K_1 : \mathbb{N} \quad K_1 : \mathbb{N} \vdash K_1 : \mathbb{N}}{K_1 : \mathbb{N} \vdash \text{plus}(K_1, K_1) : \mathbb{N}}}{K_1 : \mathbb{N}, K_2 : \mathbb{N} \vdash \text{plus}(K_1, K_2) : \mathbb{N}}}$$

$$K_1 : B, K_2 : \mathbb{N} \vdash \text{if}(K_1, K_2, \text{plus}(K_2, K_2)) : \mathbb{N}$$

dependent type theory

1960年代後半

systematic

数学的議論

AUTOMATH project

→ formal

computer check

1970年代

Martin-Löf

排中律 A V T A 背理法

選択公理 (axiom of choice)

X mechanical checking of mathematical arguments

○ formulation of a functional language.

(KENKEN) for constructive mathematics.

$$n : \mathbb{N} \vdash \text{Nat}(n) : \text{Type}$$

$$n : \mathbb{N} \vdash \text{NatList}(n) : \text{Type}$$

$$\in \text{NatList} \quad n \times m$$

$$X^n = X \times \dots \times X$$

$$n, m : \mathbb{N}$$

$$(X_i)_{i \in I}$$

添字

$$i : I \vdash X_i : \text{Set}$$

context

$$P = K_1 : A_1, K_2 : A_2, K_3 : A_3, \dots, K_n : A_n$$

$$n : \mathbb{N}, \ell : \text{NatList}(n) \quad \text{well-formed}$$

$$h : \mathbb{N}, z : \text{Matrix}(n, m) \quad \times$$

judgement

$$a : A \text{ 空間} \quad \text{"a is an object of type A"}$$

$$a \equiv b : A \quad \text{"a and b are definitionally equal objects of type A"}$$

$$a =_A b \quad \text{propositional equality}$$