

On the Achievable Rate Region in the Optimistic Sense for Separate Coding of Two Correlated General Sources*

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SUMMARY This paper is concerned with coding theorems in the optimistic sense for separate coding of two correlated general sources X_1 and X_2 . We investigate the achievable rate region $\mathcal{R}_{\text{opt}}(X_1, X_2)$ such that the decoding error probability caused by two encoders and one decoder can be arbitrarily small infinitely often under a certain rate constraint. We give an inner and an outer bounds of $\mathcal{R}_{\text{opt}}(X_1, X_2)$, where the outer bound is described by using new information-theoretic quantities. We also give two simple sufficient conditions under which the inner bound coincides with the outer bound.

key words: correlated sources, general source, optimistic coding, achievable rate region, information-spectrum methods

1. Introduction

In information-spectrum methods that originate from [3] and are described in detail in [4], fixed-length coding of a general source $X = \{X^n\}_{n=1}^\infty$ is one of fundamental problems, where, letting \mathcal{X} be a finite or countably infinite alphabet, X^n is defined as a random variable taking values in \mathcal{X}^n . The class of general sources includes vast classes of sources such as memoryless sources, Markov sources, stationary ergodic sources and stationary sources. Therefore, coding theorems for the class of general sources are valid to such classes of sources.

In coding of a general source we can formulate a new kind of coding problem in the *optimistic* sense. We are interested in the infimum of the rate R such that there exists a sequence $\{(\varphi_n, \psi_n)\}_{n=1}^\infty$ of pairs of an encoder φ_n and a decoder ψ_n satisfying for any $\gamma > 0$ both $\frac{1}{n} \log M_n \leq R + \gamma$ and $\varepsilon_n \leq \gamma$ with a subsequence $n = n_i, i \geq 1$, where M_n and ε_n denote the number of codewords and the decoding error probability, respectively. This problem was first formulated by Vembu, Verdú and Steinberg [11] and was discussed by Chen and Alajaji [1]. Results related to the optimistic coding problem can be found in [5], [9]. Quite recently, the author defined new information-theoretic quantities and clarified their operational meanings in the optimistic sense [6].

In this paper we discuss the achievable rate region in the optimistic sense for separate coding of a pair $(X_1, X_2) = \{(X_1^n, X_2^n)\}_{n=1}^\infty$ of two correlated general sources, where for each $k = 1, 2$ $X_k^n \in \mathcal{X}_k^n$ and \mathcal{X}_k denotes a finite or countably infinite alphabet. We consider two encoders $\varphi_n^{(k)}$:

$\mathcal{X}_k^n \rightarrow \{1, 2, \dots, M_n^{(k)}\}, k = 1, 2$, for coding of X_1^n and X_2^n , respectively, and a decoder ψ_n that outputs an estimate $(\hat{X}_1^n, \hat{X}_2^n) \in \mathcal{X}_1^n \times \mathcal{X}_2^n$ of (X_1^n, X_2^n) from two transmitted codewords $I_n^{(1)}$ and $I_n^{(2)}$ (see Fig. 1). We are interested in the achievable rate region $\mathcal{R}_{\text{opt}}(X_1, X_2)$ in the optimistic sense, which is defined as the collection of all the rate pairs (R_1, R_2) such that there exists a sequence $\{(\varphi_n^{(1)}, \varphi_n^{(2)}, \psi_n)\}_{n=1}^\infty$ satisfying $\frac{1}{n} \log M_n^{(k)} \leq R_k + \gamma$ ($k = 1, 2$) and $\varepsilon_n \leq \gamma$ for a subsequence $n = n_i, i \geq 1$, where $\gamma > 0$ is an arbitrarily small constant and ε_n denotes the decoding error probability.

We give an inner and an outer bounds of $\mathcal{R}_{\text{opt}}(X_1, X_2)$, where the outer bound can be expressed by using new quantities introduced in [6]. In particular, it is shown that the outer bound is easily treated by characterizing the unachievable region in the strong sense [6]. We also give two simple sufficient conditions under which the outer bound coincides with the inner bound.

The problem of separate coding of correlated sources was first formulated by Slepian and Wolf [10] for the memoryless case. A simple proof using the bin coding is given in Cover [2]. The achievable rate region for general sources X_1 and X_2 in the ordinary sense is given in Miyake and Kanaya [8]. The achievable rate region in [8] is extended to the case of ε -error by Han [4]. However, to the author's knowledge, no coding theorem in the optimistic sense is obtained so far. Our approach using the inner and the outer bounds clarifies a difference between coding in the ordinary sense and in the optimistic sense.

2. Problem Setup

In this section we consider separate coding of two correlated sources [10]. The block diagram is given in Fig. 1. Let $(X_1, X_2) = \{(X_1^n, X_2^n)\}_{n=1}^\infty$ be two correlated general sources. For each $n \geq 1$ X_1^n and X_2^n take values in \mathcal{X}_1^n and \mathcal{X}_2^n , respectively, where \mathcal{X}_1 and \mathcal{X}_2 denote finite or countably infinite alphabets. Denote by $P_{X_1^n X_2^n}$ and $P_{X_k^n | X_l^n}$ the joint probabil-

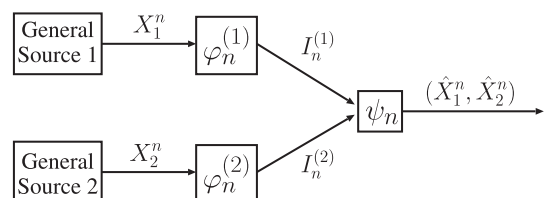


Fig. 1 Separate coding of two correlated general sources.

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ity distribution of (X_1^n, X_2^n) and the conditional probability distribution of X_k^n given X_l^n for $(k, l) = (1, 2)$ and $(2, 1)$. For each $k = 1, 2$ we define an encoder k as a deterministic mapping $\varphi_n^{(k)} : \mathcal{X}_k^n \rightarrow \mathcal{M}_n^{(k)} \stackrel{\text{def}}{=} \{1, 2, \dots, M_n^{(k)}\}$. In addition, we define a decoder as a deterministic mapping $\psi_n : \mathcal{M}_n^{(1)} \times \mathcal{M}_n^{(2)} \rightarrow \mathcal{X}_1^n \times \mathcal{X}_2^n$. The decoding error probability caused by the triplet $(\varphi_n^{(1)}, \varphi_n^{(2)}, \psi_n)$ is defined by

$$\varepsilon_n = \Pr\{\psi_n(\varphi_n^{(1)}(X_1^n), \varphi_n^{(2)}(X_2^n)) \neq (X_1^n, X_2^n)\},$$

where throughout this paper $\Pr\{\cdot\}$ denotes the probability with respect to the joint probability $P_{X_1^n X_2^n}$. Hereinafter, we call a sequence $\{(\varphi_n^{(1)}, \varphi_n^{(2)}, \psi_n)\}_{n=1}^\infty$ of the triplets of two encoders $\varphi_n^{(1)}$ and $\varphi_n^{(2)}$ and a decoder ψ_n a *code* for short.

Miyake and Kanaya [8] obtained the achievable rate region under the requirement of $\varepsilon_n \rightarrow 0$ as $n \rightarrow \infty$.

Definition 1: A rate pair (R_1, R_2) is called achievable in the *ordinary* sense if there exists a code $\{(\varphi_n^{(1)}, \varphi_n^{(2)}, \psi_n)\}_{n=1}^\infty$ satisfying

$$\limsup_{n \rightarrow \infty} \frac{1}{n} \log M_n^{(k)} \leq R_k, \quad k = 1, 2 \quad (1)$$

and

$$\lim_{n \rightarrow \infty} \varepsilon_n = 0, \quad (2)$$

where $\log(\cdot) = \log_2(\cdot)$ throughout this paper. Define the achievable rate region in the ordinary sense by

$$\begin{aligned} \mathcal{R}(X_1, X_2) = \text{Cl}(\{(R_1, R_2) : \\ (R_1, R_2) \text{ is achievable in the ordinary sense}\}), \end{aligned}$$

where $\text{Cl}(\cdot)$ denotes the closure of the set.

Theorem 1 ([8]): It holds that

$$\mathcal{R}(X_1, X_2) = \mathcal{S}(X_1, X_2),$$

where

$$\begin{aligned} \mathcal{S}(X_1, X_2) = \{(R_1, R_2) : R_1 \geq \overline{H}(X_1|X_2), \\ R_2 \geq \overline{H}(X_2|X_1) \text{ and } R_1 + R_2 \geq \overline{H}(X_1, X_2)\} \end{aligned}$$

and

$$\begin{aligned} \overline{H}(X_1, X_2) = \inf\left\{\alpha : \right. \\ \left. \liminf_{n \rightarrow \infty} \Pr\left\{\frac{1}{n} \log \frac{1}{P_{X_1^n X_2^n}(X_1^n, X_2^n)} \leq \alpha\right\} = 1\right\}, \\ \overline{H}(X_k|X_l) = \inf\left\{\alpha : \right. \\ \left. \liminf_{n \rightarrow \infty} \Pr\left\{\frac{1}{n} \log \frac{1}{P_{X_k^n|X_l^n}(X_k^n|X_l^n)} \leq \alpha\right\} = 1\right\} \end{aligned}$$

for $(k, l) = (1, 2)$ and $(2, 1)$.

It is important to notice that (1) and (2) require that for any $\gamma > 0$ $\frac{1}{n} \log M_n^{(k)} \leq R_k + \gamma$ ($k = 1, 2$) and $\varepsilon_n \leq \gamma$ for all

sufficiently large n .

In this paper, we are interested in the achievable rate region in the optimistic sense defined as follows:

Definition 2: A rate pair (R_1, R_2) is called achievable in the *optimistic* sense if there exists a code $\{(\varphi_n^{(1)}, \varphi_n^{(2)}, \psi_n)\}_{n=1}^\infty$ such that for any $\gamma > 0$ there exists a subsequence $\{n_i\}_{i=1}^\infty$ satisfying

$$\frac{1}{n} \log M_n^{(k)} \leq R_k + \gamma, \quad k = 1, 2 \quad (3)$$

and

$$\varepsilon_n \leq \gamma \quad (4)$$

for all $n = n_i$, $i \geq 1$. Define the achievable rate region in the optimistic sense by

$$\begin{aligned} \mathcal{R}_{\text{opt}}(X_1, X_2) = \text{Cl}(\{(R_1, R_2) : \\ (R_1, R_2) \text{ is achievable in the optimistic sense}\}). \end{aligned}$$

Note that in Definition 2 (3) and (4) require that for any $\gamma > 0$ $\frac{1}{n} \log M_n^{(k)} \leq R_k + \gamma$ ($k = 1, 2$) and $\varepsilon_n \leq \gamma$ simultaneously for infinitely many n . Hence, it is obvious that $\mathcal{R}(X_1, X_2) \subset \mathcal{R}_{\text{opt}}(X_1, X_2)$. While the formula of the infimum achievable rate in the optimistic sense for fixed-to-fixed length coding of a general source X is given in [1], no result is known on the achievable rate region $\mathcal{R}_{\text{opt}}(X_1, X_2)$ in the optimistic sense.

We also consider the following region.

Definition 3: A rate pair (R_1, R_2) is called *unachievable in the strong sense* if for any code $\{(\varphi_n^{(1)}, \varphi_n^{(2)}, \psi_n)\}_{n=1}^\infty$ satisfying (1) it holds that

$$\liminf_{n \rightarrow \infty} \varepsilon_n > 0. \quad (5)$$

Define the unachievable rate region in the strong sense by

$$\begin{aligned} \mathcal{U}(X_1, X_2) = \{(R_1, R_2) : \\ (R_1, R_2) \text{ is unachievable in the strong sense}\}. \end{aligned}$$

Notice that (5) means that ε_n is positive for all sufficiently large n . This $\mathcal{U}(X_1, X_2)$ is closely related to $\mathcal{R}_{\text{opt}}(X_1, X_2)$ and facilitates characterization of $\mathcal{R}_{\text{opt}}(X_1, X_2)$.

The following proposition is easily obtained from the definitions of $\mathcal{R}_{\text{opt}}(X_1, X_2)$ and $\mathcal{U}(X_1, X_2)$.

Proposition 1: Define $\mathcal{U}^\dagger(X_1, X_2) = \text{Cl}(\mathcal{U}^c(X_1, X_2))$, where the superscript c denotes the complement. Then, it holds that $\mathcal{U}^\dagger(X_1, X_2) \subset \mathcal{R}_{\text{opt}}(X_1, X_2)$.

Proof: If $(R_1, R_2) \in \mathcal{U}^c(X_1, X_2)$, there exists a code $\{(\varphi_n^{(1)}, \varphi_n^{(2)}, \psi_n)\}_{n=1}^\infty$ satisfying (1) and $\liminf_{n \rightarrow \infty} \varepsilon_n = 0$. Such a code $\{(\varphi_n^{(1)}, \varphi_n^{(2)}, \psi_n)\}_{n=1}^\infty$ satisfies, letting $\gamma > 0$ be an arbitrary constant, $\frac{1}{n} \log M_n^{(k)} \leq R_k + \gamma$, $k = 1, 2$, for all sufficiently large n and $\varepsilon_n \leq \gamma$ infinitely often. This means $\mathcal{U}^c(X_1, X_2) \subset \mathcal{R}_{\text{opt}}(X_1, X_2)$ because $(R_1, R_2) \in \mathcal{U}^c(X_1, X_2)$ is arbitrary. Since $\mathcal{R}_{\text{opt}}(X_1, X_2)$ is defined as the closure, the claim of this proposition follows. \square

Remark: In [4], [8] the achievable rate region $\mathcal{R}(X_1, X_2)$ is defined without taking the closure. In this paper, however, $\mathcal{R}(X_1, X_2)$ and $\mathcal{R}_{\text{opt}}(X_1, X_2)$ are defined as the closures so that we can ensure technical soundness especially of Proposition 1 and discussions in Sect. 4.1.

3. Main Results

Before giving inner and outer bounds of $\mathcal{R}_{\text{opt}}(X_1, X_2)$ and $\mathcal{U}(X_1, X_2)$, we define the three events for $R \geq 0$ as follows:

$$\begin{aligned} E_n^{(1)}(R) &= \left\{ (x_1^n, x_2^n) \in \mathcal{X}_1^n \times \mathcal{X}_2^n : \frac{1}{n} \log \frac{1}{P_{X_1^n|X_2^n}(x_1^n|x_2^n)} \geq R \right\}, \\ E_n^{(2)}(R) &= \left\{ (x_1^n, x_2^n) \in \mathcal{X}_1^n \times \mathcal{X}_2^n : \frac{1}{n} \log \frac{1}{P_{X_2^n|X_1^n}(x_2^n|x_1^n)} \geq R \right\}, \\ E_n^{(3)}(R) &= \left\{ (x_1^n, x_2^n) \in \mathcal{X}_1^n \times \mathcal{X}_2^n : \frac{1}{n} \log \frac{1}{P_{X_1^n X_2^n}(x_1^n, x_2^n)} \geq R \right\}. \end{aligned}$$

Furthermore, define

$$\begin{aligned} \mathcal{A}(X_1, X_2) &= \{(R_1, R_2) : \text{for any } \gamma > 0 \\ &\liminf_{n \rightarrow \infty} \Pr\{E_n^{(1)}(R_1 + \gamma) \cup E_n^{(2)}(R_2 + \gamma) \\ &\cup E_n^{(3)}(R_1 + R_2 + \gamma)\} = 0 \text{ holds}\} \end{aligned} \quad (6)$$

and

$$\begin{aligned} \mathcal{B}(X_1, X_2) &= \{(R_1, R_2) : \text{for any } \gamma > 0 \text{ all of} \\ &\liminf_{n \rightarrow \infty} \Pr\{E_n^{(1)}(R_1 + \gamma)\} = 0, \\ &\liminf_{n \rightarrow \infty} \Pr\{E_n^{(2)}(R_2 + \gamma)\} = 0, \\ &\liminf_{n \rightarrow \infty} \Pr\{E_n^{(3)}(R_1 + R_2 + \gamma)\} = 0 \text{ hold}\}. \end{aligned} \quad (7)$$

In fact, $\mathcal{B}(X_1, X_2)$ can be written in the following form.

Proposition 2: $\mathcal{B}(X_1, X_2) = \mathcal{S}^*(X_1, X_2)$, where

$$\begin{aligned} \mathcal{S}^*(X_1, X_2) &= \{(R_1, R_2) : R_1 \geq \bar{H}^*(X_1|X_2), \\ &R_2 \geq \bar{H}^*(X_2|X_1) \text{ and } R_1 + R_2 \geq \bar{H}^*(X_1, X_2)\} \end{aligned}$$

and

$$\begin{aligned} \bar{H}^*(X_1, X_2) &= \inf \left\{ \alpha : \right. \\ &\limsup_{n \rightarrow \infty} \Pr \left\{ \frac{1}{n} \log \frac{1}{P_{X_1^n|X_2^n}(X_1^n|X_2^n)} \leq \alpha \right\} = 1 \Big\}, \\ \bar{H}^*(X_k|X_l) &= \inf \left\{ \alpha : \right. \\ &\limsup_{n \rightarrow \infty} \Pr \left\{ \frac{1}{n} \log \frac{1}{P_{X_k^n|X_l^n}(X_k^n|X_l^n)} \leq \alpha \right\} = 1 \Big\} \end{aligned}$$

for $(k, l) = (1, 2)$ and $(2, 1)$.

Proof: We first prove $\mathcal{B}(X_1, X_2) \subset \mathcal{S}^*(X_1, X_2)$. Let (R_1, R_2) be an arbitrary element of $\mathcal{B}(X_1, X_2)$. Then, letting $\gamma > 0$ be an arbitrary constant, we have

$$\liminf_{n \rightarrow \infty} \Pr\{E_n^{(1)}(R_1 + \gamma)\}$$

$$= \liminf_{n \rightarrow \infty} \Pr \left\{ \frac{1}{n} \log \frac{1}{P_{X_1^n|X_2^n}(X_1^n|X_2^n)} \geq R_1 + \gamma \right\} = 0,$$

which implies that

$$\limsup_{n \rightarrow \infty} \Pr \left\{ \frac{1}{n} \log \frac{1}{P_{X_1^n|X_2^n}(X_1^n|X_2^n)} \leq R_1 + \gamma \right\} = 1 \quad (8)$$

for any $\gamma > 0$. Note that we can use \leq instead of $<$ in (8). Hence, it holds that $\bar{H}^*(X_1|X_2) \leq R_1$ owing to the definition of $\bar{H}^*(X_1|X_2)$. Similarly, we can obtain $\bar{H}^*(X_2|X_1) \leq R_2$ and $\bar{H}^*(X_1, X_2) \leq R_1 + R_2$. Thus, we obtain $(R_1, R_2) \in \mathcal{S}^*(X_1, X_2)$. Since $(R_1, R_2) \in \mathcal{B}(X_1, X_2)$ is arbitrary, $\mathcal{B}(X_1, X_2) \subset \mathcal{S}^*(X_1, X_2)$ is established.

Next, we prove $\mathcal{S}^*(X_1, X_2) \subset \mathcal{B}(X_1, X_2)$. Fix $(R_1, R_2) \in \mathcal{S}^*(X_1, X_2)$ arbitrarily. Since $R_1 \geq \bar{H}^*(X_1|X_2)$ is guaranteed, for any $\gamma > 0$ it follows that

$$\begin{aligned} \Pr\{E_n^{(1)}(R_1 + \gamma)\} &= \Pr \left\{ \frac{1}{n} \log \frac{1}{P_{X_1^n|X_2^n}(X_1^n|X_2^n)} \geq R_1 + \gamma \right\} \\ &\leq \Pr \left\{ \frac{1}{n} \log \frac{1}{P_{X_1^n|X_2^n}(X_1^n|X_2^n)} > \bar{H}^*(X_1|X_2) + \frac{\gamma}{2} \right\} \\ &= 1 - \Pr \left\{ \frac{1}{n} \log \frac{1}{P_{X_1^n|X_2^n}(X_1^n|X_2^n)} \leq \bar{H}^*(X_1|X_2) + \frac{\gamma}{2} \right\}. \end{aligned}$$

By taking the limit inferior of both sides, we have

$$\begin{aligned} \liminf_{n \rightarrow \infty} \Pr\{E_n^{(1)}(R_1 + \gamma)\} \\ \leq 1 - \limsup_{n \rightarrow \infty} \Pr \left\{ \frac{1}{n} \log \frac{1}{P_{X_1^n|X_2^n}(X_1^n|X_2^n)} \leq \bar{H}^*(X_1|X_2) + \frac{\gamma}{2} \right\} \\ = 0 \end{aligned}$$

due to the definition of $\bar{H}^*(X_1|X_2)$. Since $\Pr\{E_n^{(1)}(R_1 + \gamma)\} \geq 0$ for all $n \geq 1$, this establishes $\liminf_{n \rightarrow \infty} \Pr\{E_n^{(1)}(R_1 + \gamma)\} = 0$. Similarly, we can prove $\liminf_{n \rightarrow \infty} \Pr\{E_n^{(2)}(R_2 + \gamma)\} = 0$ and $\liminf_{n \rightarrow \infty} \Pr\{E_n^{(3)}(R_1 + R_2 + \gamma)\} = 0$. Thus, we obtain $(R_1, R_2) \in \mathcal{B}(X_1, X_2)$. Since $(R_1, R_2) \in \mathcal{S}^*(X_1, X_2)$ is arbitrary, $\mathcal{S}^*(X_1, X_2) \subset \mathcal{B}(X_1, X_2)$ follows. \square

Proposition 3: $\mathcal{A}(X_1, X_2) \subset \mathcal{B}(X_1, X_2)$.

Proof: For any $(R_1, R_2) \in \mathcal{A}(X_1, X_2)$ it holds that

$$\begin{aligned} \liminf_{n \rightarrow \infty} \Pr\{E_n^{(1)}(R_1 + \gamma)\} &= \liminf_{n \rightarrow \infty} \Pr\{E_n^{(2)}(R_2 + \gamma)\} \\ &= \liminf_{n \rightarrow \infty} \Pr\{E_n^{(3)}(R_1 + R_2 + \gamma)\} = 0 \end{aligned}$$

because all of $E_n^{(1)}(R_1 + \gamma)$, $E_n^{(2)}(R_2 + \gamma)$ and $E_n^{(3)}(R_1 + R_2 + \gamma)$ are subsets of $E_n^{(1)}(R_1 + \gamma) \cup E_n^{(2)}(R_2 + \gamma) \cup E_n^{(3)}(R_1 + R_2 + \gamma)$. Thus, the claim of this proposition follows immediately. \square

Now, we are ready to state the main results.

Theorem 2: For any two correlated general sources X_1 and X_2 we have

$$\mathcal{A}(X_1, X_2) \subset \mathcal{U}^\dagger(X_1, X_2) \subset \mathcal{R}_{\text{opt}}(X_1, X_2) \subset \mathcal{B}(X_1, X_2).$$

We use the following two lemmas in the proof of Theorem 2.

Lemma 1 ([4]): Let $M_n^{(1)}$ and $M_n^{(2)}$ be positive integers arbitrarily given. Then, there exists a code $\{(\varphi_n^{(1)}, \varphi_n^{(2)}, \psi_n)\}_{n=1}^\infty$ satisfying

$$\begin{aligned} \varepsilon_n \leq & \Pr\left\{\frac{1}{n} \log \frac{1}{P_{X_1^n|X_2^n}(X_1^n|X_2^n)} \geq \frac{1}{n} \log M_n^{(1)} - \gamma \right. \\ & \text{or } \frac{1}{n} \log \frac{1}{P_{X_2^n|X_1^n}(X_2^n|X_1^n)} \geq \frac{1}{n} \log M_n^{(2)} - \gamma \\ & \left. \text{or } \frac{1}{n} \log \frac{1}{P_{X_1^n X_2^n}(X_1^n X_2^n)} \geq \frac{1}{n} \log M_n^{(1)} M_n^{(2)} - \gamma \right\} \\ & + 3 \cdot 2^{-n\gamma} \end{aligned} \quad (9)$$

for any $\gamma > 0$ and $n \geq 1$.

Lemma 2 ([4]): Any code $\{(\varphi_n^{(1)}, \varphi_n^{(2)}, \psi_n)\}_{n=1}^\infty$ satisfies

$$\begin{aligned} \varepsilon_n \geq & \Pr\left\{\frac{1}{n} \log \frac{1}{P_{X_1^n|X_2^n}(X_1^n|X_2^n)} \geq \frac{1}{n} \log M_n^{(1)} + \gamma \right. \\ & \text{or } \frac{1}{n} \log \frac{1}{P_{X_2^n|X_1^n}(X_2^n|X_1^n)} \geq \frac{1}{n} \log M_n^{(2)} + \gamma \\ & \left. \text{or } \frac{1}{n} \log \frac{1}{P_{X_1^n X_2^n}(X_1^n X_2^n)} \geq \frac{1}{n} \log M_n^{(1)} M_n^{(2)} + \gamma \right\} \\ & - 3 \cdot 2^{-n\gamma} \end{aligned} \quad (10)$$

for all $n \geq 1$ and $\gamma > 0$.

Proof of Theorem 2: Since Proposition 1 guarantees $\mathcal{U}^\dagger(X_1, X_2) \subset \mathcal{R}_{\text{opt}}(X_1, X_2)$, we have only to prove $\mathcal{A}(X_1, X_2) \subset \mathcal{U}^\dagger(X_1, X_2)$ and $\mathcal{R}_{\text{opt}}(X_1, X_2) \subset \mathcal{B}(X_1, X_2)$.

First, we prove $\mathcal{A}(X_1, X_2) \subset \mathcal{U}^\dagger(X_1, X_2)$. We use the idea in [4, Theorem 7.4.1]. Fix $\gamma > 0$ arbitrarily. We first show that $(R_1, R_2) \in \mathcal{A}(X_1, X_2)$ implies $(R_1 + 2\gamma, R_2 + 2\gamma) \in \mathcal{U}^c(X_1, X_2)$. Set $M_n^{(1)} = 2^{n(R_1 + 2\gamma)}$ and $M_n^{(2)} = 2^{n(R_2 + 2\gamma)}$. Clearly, these $M_n^{(1)}$ and $M_n^{(2)}$ satisfy

$$\limsup_{n \rightarrow \infty} \frac{1}{n} \log M_n^{(k)} \leq R_k + 2\gamma, \quad k = 1, 2. \quad (11)$$

Since Lemma 1 guarantees the existence of a code $\{(\varphi_n^{(1)}, \varphi_n^{(2)}, \psi_n)\}_{n=1}^\infty$ satisfying (9), we have

$$\begin{aligned} \varepsilon_n \leq & \Pr\left\{\frac{1}{n} \log \frac{1}{P_{X_1^n|X_2^n}(X_1^n|X_2^n)} \geq R_1 + \gamma \right. \\ & \text{or } \frac{1}{n} \log \frac{1}{P_{X_2^n|X_1^n}(X_2^n|X_1^n)} \geq R_2 + \gamma \\ & \left. \text{or } \frac{1}{n} \log \frac{1}{P_{X_1^n X_2^n}(X_1^n X_2^n)} \geq R_1 + R_2 + 3\gamma \right\} \\ & + 3 \cdot 2^{-n\gamma} \\ \leq & \Pr\{E_n^{(1)}(R_1 + \gamma) \cup E_n^{(2)}(R_2 + \gamma) \\ & \cup E_n^{(3)}(R_1 + R_2 + \gamma)\} + 3 \cdot 2^{-n\gamma}. \end{aligned}$$

Since $(R_1, R_2) \in \mathcal{A}(X_1, X_2)$ by assumption, by taking the limit inferior of both sides we have

$$\liminf_{n \rightarrow \infty} \varepsilon_n = 0,$$

which, together with (11), implies that $(R_1 + 2\gamma, R_2 + 2\gamma) \in$

$\mathcal{U}^c(X_1, X_2)$. Since $\gamma > 0$ is arbitrary, by letting $\gamma \downarrow 0$ we can conclude that $(R_1, R_2) \in \mathcal{U}^\dagger(X_1, X_2)$. This completes the proof of $\mathcal{A}(X_1, X_2) \subset \mathcal{U}^\dagger(X_1, X_2)$.

Next, we prove $\mathcal{R}_{\text{opt}}(X_1, X_2) \subset \mathcal{B}(X_1, X_2)$. In view of Proposition 2, it suffices to prove that any $(R_1, R_2) \in \mathcal{R}_{\text{opt}}(X_1, X_2)$ belongs to $\mathcal{S}^*(X_1, X_2)$, i.e., all of (a) $R_1 \geq \overline{H}^*(X_1|X_2)$, (b) $R_2 \geq \overline{H}^*(X_2|X_1)$, and (c) $R_1 + R_2 \geq \overline{H}^*(X_1, X_1)$ are satisfied.

Hereinafter, we prove $R_1 \geq \overline{H}^*(X_1|X_2)$ by contradiction. Assume that $R_1 < \overline{H}^*(X_1|X_2)$ is satisfied. Then, there exists a constant $\gamma_0 > 0$ such that $R_1 \leq \overline{H}^*(X_1|X_2) - 3\gamma_0$. In addition, we should note that $\overline{H}^*(X_1|X_2)$ can be expressed in the following form:

$$\begin{aligned} \overline{H}^*(X_1|X_2) = & \sup\{\beta : \\ & \liminf_{n \rightarrow \infty} \Pr\left\{\frac{1}{n} \log \frac{1}{P_{X_1^n|X_2^n}(X_1^n|X_2^n)} \geq \beta\right\} > 0\} \end{aligned}$$

(see [6, Appendix]). Therefore, there exist a constant $\delta_0 > 0$ and an integer n_0 such that

$$\begin{aligned} \Pr\left\{\frac{1}{n} \log \frac{1}{P_{X_1^n|X_2^n}(X_1^n|X_2^n)} \geq \overline{H}^*(X_1|X_2) - \gamma_0\right\} \\ \geq 2\delta_0 \quad \text{for all } n \geq n_0. \end{aligned} \quad (12)$$

Furthermore, Lemma 2 tells us that for all $n \geq 1$ any code $\{(\varphi_n^{(1)}, \varphi_n^{(2)}, \psi_n)\}_{n=1}^\infty$ satisfies (10) with $\gamma = \gamma_0$, which leads to

$$\begin{aligned} \varepsilon_n \geq & \Pr\left\{\frac{1}{n} \log \frac{1}{P_{X_1^n|X_2^n}(X_1^n|X_2^n)} \geq \frac{1}{n} \log M_n^{(1)} + \gamma_0\right\} \\ & - 3 \cdot 2^{-n\gamma_0} \quad \text{for all } n \geq 1. \end{aligned} \quad (13)$$

Now, assume that (R_1, R_2) is achievable in the optimistic sense. Then, for any $\gamma \in (0, \min\{\gamma_0, \delta_0\})$ there exists a code $\{(\varphi_n^{(1)}, \varphi_n^{(2)}, \psi_n)\}_{n=1}^\infty$ with a subsequence $\{n_i\}_{i=1}^\infty$ satisfying $\frac{1}{n} \log M_n^{(1)} \leq R_1 + \gamma \leq \overline{H}^*(X_1|X_2) - 3\gamma_0 + \gamma \leq \overline{H}^*(X_1|X_2) - 2\gamma_0$ and $\varepsilon_n \leq \gamma$ for all $n = n_i, i \geq 1$. Then, (13) leads to

$$\begin{aligned} \varepsilon_n \geq & \Pr\left\{\frac{1}{n} \log \frac{1}{P_{X_1^n|X_2^n}(X_1^n|X_2^n)} \geq \overline{H}^*(X_1|X_2) - \gamma_0\right\} \\ & - 3 \cdot 2^{-n\gamma_0} \quad \text{for all } n = n_i, i \geq 1, \end{aligned} \quad (14)$$

where the left side is less than γ for all $n = n_i, i \geq 1$, and the right side is greater than $\delta_0 (> \gamma)$ for all sufficiently large $n = n_i$ due to (12). This is a contradiction.

Since this argument can be applied to establishing (b) $R_2 \geq \overline{H}^*(X_2|X_1)$ and (c) $R_1 + R_2 \geq \overline{H}^*(X_1, X_1)$ as well, we can conclude that $(R_1, R_2) \in \mathcal{S}^*(X_1, X_2)$ provided that (R_1, R_2) is achievable in the optimistic sense. This establishes $\mathcal{R}_{\text{opt}}(X_1, X_2) \subset \mathcal{S}^*(X_1, X_2)$ because $\mathcal{S}^*(X_1, X_2)$ includes the boundaries. \square

4. Discussions

4.1 Characterization of $\mathcal{R}(X_1, X_2)$

So far, we have established an inner and an outer bounds

of $\mathcal{R}_{\text{opt}}(X_1, X_2)$. In this subsection we revisit $\mathcal{R}(X_1, X_2)$ in Definition 1 and characterize $\mathcal{R}(X_1, X_2)$ in the same manner as Theorem 2.

To this end, define $\tilde{\mathcal{A}}(X_1, X_2)$ and $\tilde{\mathcal{B}}(X_1, X_2)$ by replacing the limit inferiors in (6) and (7) with the limit superiors, respectively. Define $\mathcal{W}(X_1, X_2)$ by replacing the limit inferior in (5) with the limit superior. Then, it is easily verified that $\mathcal{R}(X_1, X_2)$ satisfies

$$\tilde{\mathcal{A}}(X_1, X_2) \subset \mathcal{W}^\dagger(X_1, X_2) \subset \mathcal{R}(X_1, X_2) \subset \tilde{\mathcal{B}}(X_1, X_2)$$

similarly to the proof of Theorem 2, where $\mathcal{W}^\dagger(X_1, X_2)$ denotes the closure of the complement of $\mathcal{W}(X_1, X_2)$. We can also verify that $\tilde{\mathcal{B}}(X_1, X_2)$ is expressed as $\mathcal{S}(X_1, X_2)$ in Theorem 1 similarly to Proposition 2.

In the ordinary case, however, we can prove $\tilde{\mathcal{B}}(X_1, X_2) \subset \tilde{\mathcal{A}}(X_1, X_2)$ as well and therefore $\mathcal{W}^\dagger(X_1, X_2) = \mathcal{R}(X_1, X_2) = \mathcal{S}(X_1, X_2)$ without any assumption on (X_1, X_2) . This explains a reason why $\mathcal{R}(X_1, X_2)$ is expressed in a closed form and is coincident with $\mathcal{S}(X_1, X_2)$.

In order to verify $\tilde{\mathcal{B}}(X_1, X_2) \subset \tilde{\mathcal{A}}(X_1, X_2)$, let (R_1, R_2) be an arbitrary element of $\tilde{\mathcal{B}}(X_1, X_2)$. By the definition of $\tilde{\mathcal{B}}(X_1, X_2)$, it holds that $\limsup_{n \rightarrow \infty} \Pr\{E_n^{(k)}(R_k + \gamma)\} = 0$ for $k = 1, 2$ and $\limsup_{n \rightarrow \infty} \Pr\{E_n^{(3)}(R_1 + R_2 + \gamma)\} = 0$. Then, we have

$$\begin{aligned} & \limsup_{n \rightarrow \infty} \Pr\{E_n^{(1)}(R_1 + \gamma) \cup E_n^{(2)}(R_2 + \gamma) \cup E_n^{(3)}(R_1 + R_2 + \gamma)\} \\ & \leq \limsup_{n \rightarrow \infty} \Pr\{E_n^{(1)}(R_1 + \gamma)\} + \limsup_{n \rightarrow \infty} \Pr\{E_n^{(2)}(R_2 + \gamma)\} \\ & \quad + \limsup_{n \rightarrow \infty} \Pr\{E_n^{(3)}(R_1 + R_2 + \gamma)\} \\ & = 0, \end{aligned} \quad (15)$$

where the inequality follows from the union bound and $\limsup_{n \rightarrow \infty} (a_n + b_n) \leq \limsup_{n \rightarrow \infty} a_n + \limsup_{n \rightarrow \infty} b_n$ for any real-valued sequences $\{a_n\}_{n=1}^\infty$ and $\{b_n\}_{n=1}^\infty$. Since the left side is nonnegative, (15) implies that $(R_1, R_2) \in \tilde{\mathcal{A}}(X_1, X_2)$. This establishes $\tilde{\mathcal{B}}(X_1, X_2) \subset \tilde{\mathcal{A}}(X_1, X_2)$.

Notice that we cannot use the same argument to establish $\mathcal{B}(X_1, X_2) \subset \mathcal{A}(X_1, X_2)$ because only the opposite side the inequality holds with respect to the limit inferior. In Example 3 in Sect. 4.2 we will see an example of (X_1, X_2) satisfying $\mathcal{A}(X_1, X_2) \subsetneq \mathcal{B}(X_1, X_2)$.

4.2 Sufficient Conditions for $\mathcal{A}(X_1, X_2) = \mathcal{B}(X_1, X_2)$

In this subsection, we investigate sufficient conditions on (X_1, X_2) under which $\mathcal{A}(X_1, X_2) = \mathcal{B}(X_1, X_2)$ holds. Theorem 2 and Proposition 2 guarantee that $\mathcal{R}_{\text{opt}}(X_1, X_2) = \mathcal{U}^\dagger(X_1, X_2) = \mathcal{S}^*(X_1, X_2)$ under such sufficient conditions.

We begin with the simple sufficient condition given in the following proposition. This sufficient condition is shown to be valid by using the argument given in Sect. 4.1.

Proposition 4: If $\overline{H}(X_1|X_2) = \overline{H}^*(X_1|X_2)$, $\overline{H}(X_2|X_1) = \overline{H}^*(X_2|X_1)$ and $\overline{H}(X_1, X_2) = \overline{H}^*(X_1, X_2)$, then it holds that $\mathcal{R}_{\text{opt}}(X_1, X_2) = \mathcal{U}^\dagger(X_1, X_2) = \mathcal{S}^*(X_1, X_2)$.

Proof: Clearly, $\mathcal{S}(X_1, X_2) = \mathcal{S}^*(X_1, X_2)$ holds by

the assumption. So far, we have already established that $\mathcal{B}(X_1, X_2) = \mathcal{S}^*(X_1, X_2)$ in Proposition 2 and $\tilde{\mathcal{B}}(X_1, X_2) = \mathcal{S}(X_1, X_2)$ in Sect. 4.1. Hence, $\mathcal{S}(X_1, X_2) = \mathcal{S}^*(X_1, X_2)$ implies $\mathcal{B}(X_1, X_2) = \tilde{\mathcal{B}}(X_1, X_2)$. In addition, we should note $\tilde{\mathcal{A}}(X_1, X_2) \subset \mathcal{A}(X_1, X_2)$, which immediately follows from their definitions. Then, in view of Theorem 2 it holds that

$$\begin{aligned} \tilde{\mathcal{A}}(X_1, X_2) & \subset \mathcal{A}(X_1, X_2) \subset \mathcal{U}^\dagger(X_1, X_2) \\ & \subset \mathcal{R}_{\text{opt}}(X_1, X_2) \subset \mathcal{B}(X_1, X_2) = \tilde{\mathcal{B}}(X_1, X_2). \end{aligned} \quad (16)$$

Since we have proved that $\tilde{\mathcal{A}}(X_1, X_2) = \tilde{\mathcal{B}}(X_1, X_2)$ in Sect. 4.1, the claim in this proposition follows. \square

Example 1: Let $P_{X_1 X_2}$ be a joint probability distribution on $X_1 \times X_2$ satisfying $H(X_1 X_2) < \infty$, where $H(X_1 X_2)$ denotes the joint entropy of X_1 and X_2 . Let X_1 and X_2 be two correlated stationary memoryless sources induced by $P_{X_1 X_2}$. Clearly, by the weak law of large numbers we have $\overline{H}(X_1|X_2) = \overline{H}^*(X_1|X_2) = H(X_1|X_2)$, $\overline{H}(X_2|X_1) = \overline{H}^*(X_2|X_1) = H(X_2|X_1)$ and $\overline{H}(X_1, X_2) = \overline{H}^*(X_1, X_2) = H(X_1 X_2)$ for these sources, where $H(X_k|X_l)$ denotes the conditional entropy of X_k given X_l for $(k, l) = (1, 2)$ and $(2, 1)$. Hence, Proposition 4 guarantees that

$$\begin{aligned} \mathcal{R}(X_1, X_2) & = \mathcal{R}_{\text{opt}}(X_1, X_2) \\ & = \{(R_1, R_2) : R_1 \geq H(X_1|X_2), R_2 \geq H(X_2|X_1) \\ & \quad \text{and } R_1 + R_2 \geq H(X_1 X_2)\}. \end{aligned}$$

Next, we introduce a new class of correlated sources.

Definition 4: Let X_1 and X_2 be two correlated general sources. We say that (X_1, X_2) is *synchronizing* if for any $\gamma > 0$ and $\eta \in (0, 1)$ there exists a subsequence $\{n_i\}_{i=1}^\infty$ such that all of

$$\Pr \left\{ \frac{1}{n} \log \frac{1}{P_{X_1^n|X_2^n}(X_1^n|X_2^n)} \geq \overline{H}^*(X_1|X_2) + \gamma \right\} \leq \eta, \quad (17)$$

$$\Pr \left\{ \frac{1}{n} \log \frac{1}{P_{X_2^n|X_1^n}(X_2^n|X_1^n)} \geq \overline{H}^*(X_2|X_1) + \gamma \right\} \leq \eta, \quad (18)$$

$$\Pr \left\{ \frac{1}{n} \log \frac{1}{P_{X_1^n X_2^n}(X_1^n, X_2^n)} \geq \overline{H}^*(X_1, X_2) + \gamma \right\} \leq \eta \quad (19)$$

are satisfied for all $n = n_i$, $i \geq 1$.

Note that owing to the definitions of $\overline{H}^*(X_1|X_2)$, $\overline{H}^*(X_2|X_1)$ and $\overline{H}^*(X_1, X_2)$, each of the above three inequalities in Definition 4 is satisfied infinitely often. The synchronizing property in Definition 4 actually requires that all the three inequalities are satisfied simultaneously with the same subsequence $\{n_i\}_{i=1}^\infty$.

The following proposition indicates that the synchronizing property of (X_1, X_2) is another sufficient condition for $\mathcal{A}(X_1, X_2) = \mathcal{B}(X_1, X_2)$.

Proposition 5: If (X_1, X_2) is synchronizing, $\mathcal{R}_{\text{opt}}(X_1, X_2) = \mathcal{U}^\dagger(X_1, X_2) = \mathcal{S}^*(X_1, X_2)$.

Proof: In view of Theorem 2, it suffices to establish

$\mathcal{B}(X_1, X_2) \subset \mathcal{A}(X_1, X_2)$. Fix $(R_1, R_2) \in \mathcal{B}(X_1, X_2)$ arbitrarily. Then, Proposition 2 tells us that all of $R_1 \geq \overline{H}^*(X_1|X_2)$, $R_2 \geq \overline{H}^*(X_2|X_1)$ and $R_1 + R_2 \geq \overline{H}^*(X_1, X_2)$ are satisfied. Since (X_1, X_2) is assumed to be synchronizing, for any constants $\gamma > 0$ and $\eta \in (0, 1)$ there exists a subsequence $\{n_i\}_{i=1}^\infty$ satisfying

$$\begin{aligned} & \Pr\{E_n^{(1)}(R_1 + \gamma)\} \\ &= \Pr\left\{\frac{1}{n} \log \frac{1}{P_{X_1^n|X_2^n}(X_1^n|X_2^n)} \geq R_1 + \gamma\right\} \\ &\leq \Pr\left\{\frac{1}{n} \log \frac{1}{P_{X_1^n|X_2^n}(X_1^n|X_2^n)} \geq \overline{H}^*(X_1|X_2) + \gamma\right\} \\ &\leq \eta \quad \text{for all } n = n_i \text{ and } i \geq 1. \end{aligned}$$

Similarly, we have $\Pr\{E_n^{(2)}(R_2 + \gamma)\} \leq \eta$ and $\Pr\{E_n^{(3)}(R_1 + R_2 + \gamma)\} \leq \eta$ for all $n = n_i$, $i \geq 1$. Therefore, it holds that

$$\begin{aligned} & \Pr\{E_n^{(1)}(R_1 + \gamma) \cup E_n^{(2)}(R_2 + \gamma) \cup E_n^{(3)}(R_1 + R_2 + \gamma)\} \\ &\leq \Pr\{E_n^{(1)}(R_1 + \gamma)\} + \Pr\{E_n^{(2)}(R_2 + \gamma)\} \\ &\quad + \Pr\{E_n^{(3)}(R_1 + R_2 + \gamma)\} \\ &\leq 3\eta \quad \text{for all } n = n_i \text{ and } i \geq 1, \end{aligned}$$

which means that the limit inferior of the left side is equal to zero. Since $(R_1, R_2) \in \mathcal{B}(X_1, X_2)$ is arbitrary, $\mathcal{B}(X_1, X_2) \subset \mathcal{A}(X_1, X_2)$ follows. \square

The notion of synchronizing property in Definition 4 leads to an interesting implication on inner bounds of $\mathcal{R}_{\text{opt}}(X_1, X_2)$. Assume that for any $\gamma > 0$ there exists a subsequence $\{n_i\}_{i=1}^\infty$ such that (17) and (19) hold for all $n = n_i$, $i \geq 1$. Then, we can show that

$$\begin{aligned} \mathcal{V}_1(X_1, X_2) &= \{(R_1, R_2) : R_1 \geq \overline{H}^*(X_1|X_2), \\ &\quad R_2 \geq \overline{H}(X_2|X_1) \text{ and } R_1 + R_2 \geq \overline{H}^*(X_1, X_2)\} \end{aligned}$$

is a subset of $\mathcal{R}_{\text{opt}}(X_1, X_2)$. This fact is verified by checking $\mathcal{V}_1(X_1, X_2) \subset \mathcal{A}(X_1, X_2)$ as follows. Fix $(R_1, R_2) \in \mathcal{V}_1(X_1, X_2)$ arbitrarily. Then, we have

$$\begin{aligned} & \Pr\{E_n^{(1)}(R_1 + \gamma)\} + \Pr\{E_n^{(3)}(R_1 + R_2 + \gamma)\} \leq 2\eta \\ &\quad \text{for all } n = n_i \text{ and } i \geq 1, \text{ and} \\ & \Pr\{E_n^{(2)}(R_2 + \gamma)\} \leq \eta \quad \text{for all sufficiently large } n. \end{aligned}$$

Thus, there exists an integer i_0 such that

$$\begin{aligned} & \Pr\{E_n^{(1)}(R_1 + \gamma) \cup E_n^{(2)}(R_2 + \gamma) \cup E_n^{(3)}(R_1 + R_2 + \gamma)\} \leq 3\eta \\ &\quad \text{for all } n = n_i \text{ and } i \geq i_0, \end{aligned}$$

which means the limit inferior of the left side is equal to zero. On the other hand, we can similarly prove that

$$\begin{aligned} \mathcal{V}_2(X_1, X_2) &= \{(R_1, R_2) : R_1 \geq \overline{H}(X_1|X_2), \\ &\quad R_2 \geq \overline{H}^*(X_2|X_1) \text{ and } R_1 + R_2 \geq \overline{H}^*(X_1, X_2)\} \end{aligned}$$

is a subset of $\mathcal{A}(X_1, X_2)$ with no assumption on (X_1, X_2) .

Example 2: Let $P_{X_1 X_2}^{(1)}$ and $P_{X_1 X_2}^{(2)}$ be the joint probability distributions on $X_1 \times X_2$ satisfying $H(X_1^{(1)}|X_2^{(1)}) \leq H(X_1^{(2)}|X_2^{(2)})$, $H(X_2^{(1)}|X_1^{(1)}) \leq H(X_2^{(2)}|X_1^{(2)})$ and $H(X_1^{(1)}X_2^{(1)}) \leq H(X_1^{(2)}X_2^{(2)}) <$

∞ . Define $(X_1, X_2) = \{(X_1^n, X_2^n)\}_{n=1}^\infty$ as the pair of general sources satisfying

$$P_{X_1^n X_2^n}(x_1^n, x_2^n) = \begin{cases} \prod_{i=1}^n P_{X_1 X_2}^{(1)}(x_i, x_i), & \text{if } n \text{ is odd,} \\ \prod_{i=1}^n P_{X_1 X_2}^{(2)}(x_i, x_i), & \text{if } n \text{ is even,} \end{cases}$$

where $x_1^n = (x_1, \dots, x_n)$ and $x_2^n = (x_1, \dots, x_n)$. Since it holds that $\overline{H}(X_1|X_2) = H(X_1^{(2)}|X_2^{(2)})$, $\overline{H}(X_2|X_1) = H(X_2^{(2)}|X_1^{(2)})$, and $\overline{H}(X_1 X_2) = H(X_1^{(2)} X_2^{(2)})$, we have

$$\begin{aligned} \mathcal{R}(X_1, X_2) &= \{(R_1, R_2) : R_1 \geq H(X_1^{(2)}|X_2^{(2)}), \\ &\quad R_2 \geq H(X_2^{(2)}|X_1^{(2)}) \text{ and } R_1 + R_2 \geq H(X_1^{(2)} X_2^{(2)})\}. \end{aligned}$$

In addition, since (X_1, X_2) is synchronizing and satisfies $\overline{H}^*(X_1|X_2) = H(X_1^{(1)}|X_2^{(1)})$, $\overline{H}^*(X_2|X_1) = H(X_2^{(1)}|X_1^{(1)})$, and $\overline{H}^*(X_1 X_2) = H(X_1^{(1)} X_2^{(1)})$, we have

$$\begin{aligned} \mathcal{R}_{\text{opt}}(X_1, X_2) &= \{(R_1, R_2) : R_1 \geq H(X_1^{(1)}|X_2^{(1)}), \\ &\quad R_2 \geq H(X_2^{(1)}|X_1^{(1)}) \text{ and } R_1 + R_2 \geq H(X_1^{(1)} X_2^{(1)})\}. \end{aligned}$$

Figure 2 shows $\mathcal{R}(X_1, X_2)$ and $\mathcal{R}_{\text{opt}}(X_1, X_2)$ in this example.

Example 3: Suppose that $P_{X_1 X_2}^{(1)}$ and $P_{X_1 X_2}^{(2)}$ satisfy $H(X_1^{(1)}|X_2^{(1)}) \leq H(X_1^{(2)}|X_2^{(2)})$, $H(X_1^{(1)} X_2^{(1)}) \leq H(X_1^{(2)} X_2^{(2)}) < \infty$, and $H(X_2^{(2)}|X_1^{(2)}) \leq H(X_2^{(1)}|X_1^{(1)})$ in Example 2. In this case, we can easily verify from Theorems 1 and 2 that

$$\begin{aligned} \mathcal{R}(X_1, X_2) &= \{(R_1, R_2) : R_1 \geq H(X_1^{(2)}|X_2^{(2)}), \\ &\quad R_2 \geq H(X_2^{(1)}|X_1^{(1)}) \text{ and } R_1 + R_2 \geq H(X_1^{(2)} X_2^{(2)})\} \end{aligned}$$

and

$$\mathcal{V}(X_1, X_2) \subset \mathcal{R}_{\text{opt}}(X_1, X_2) \subset \mathcal{S}^*(X_1, X_2),$$

where $\mathcal{V}(X_1, X_2) = \mathcal{V}_1(X_1, X_2) \cup \mathcal{V}_2(X_1, X_2)$,

$$\begin{aligned} \mathcal{V}_i(X_1, X_2) &= \{(R_1, R_2) : R_1 \geq H(X_1^{(i)}|X_2^{(i)}), \\ &\quad R_2 \geq H(X_2^{(i)}|X_1^{(i)}) \text{ and } R_1 + R_2 \geq H(X_1^{(i)} X_2^{(i)})\} \end{aligned}$$

for $i = 1, 2$ and

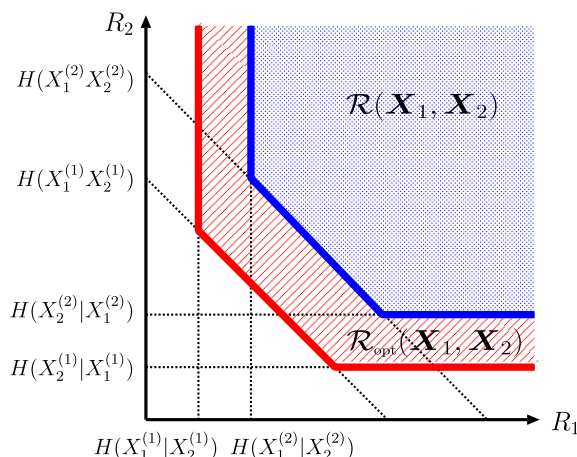


Fig. 2 $\mathcal{R}(X_1, X_2)$ and $\mathcal{R}_{\text{opt}}(X_1, X_2)$ in Example 2.

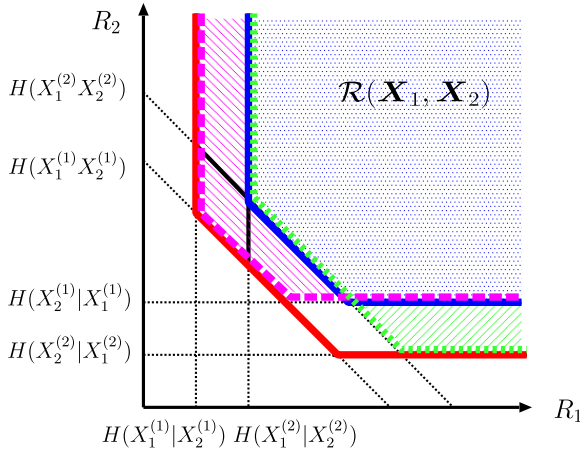


Fig. 3 $\mathcal{S}^*(X_1, X_2)$, $\mathcal{V}(X_1, X_2)$ and $\mathcal{R}(X_1, X_2)$ in Example 3. The broken (purple) line and the dotted (green) line correspond to the boundaries of $\mathcal{V}_1(X_1, X_2)$ and $\mathcal{V}_2(X_1, X_2)$, respectively.

$$\mathcal{S}^*(X_1, X_2) = \{(R_1, R_2) : R_1 \geq H(X_1^{(1)}|X_2^{(1)}), \\ R_2 \geq H(X_2^{(2)}|X_1^{(2)}) \text{ and } R_1 + R_2 \geq H(X_1^{(1)}X_2^{(1)})\}.$$

Figure 3 shows the relationship of $\mathcal{S}^*(X_1, X_2)$, $\mathcal{V}(X_1, X_2)$ and $\mathcal{R}(X_1, X_2)$ in this example.

In Example 3, we can actually prove that $\mathcal{R}_{\text{opt}}(X_1, X_2) = \mathcal{V}(X_1, X_2)$ by using a contradiction argument. Suppose that $(R_1, R_2) \in \mathcal{S}^*(X_1, X_2) \setminus \mathcal{V}(X_1, X_2)$ is achievable in the optimistic sense. Then, there exists a code $\{(\varphi_n^{(1)}, \varphi_n^{(2)}, \psi_n)\}_{n=1}^\infty$ satisfying $\frac{1}{n} \log M_n^{(k)} \leq R_k + \gamma$, $k = 1, 2$, and $\varepsilon_n \leq \gamma$ for any $\gamma > 0$ with some subsequence $\{n_i\}_{i=1}^\infty$. We first note that we can choose a sufficiently small $\gamma_0 \in (0, \frac{1}{3})$ such that $(R_1 + 2\gamma_0, R_2 + 2\gamma_0) \in \mathcal{S}^*(X_1, X_2) \setminus \mathcal{V}(X_1, X_2)$. Then, the code $\{(\varphi_n^{(1)}, \varphi_n^{(2)}, \psi_n)\}_{n=1}^\infty$ satisfies $\frac{1}{n} \log M_n^{(k)} \leq R_k + \gamma_0$, $k = 1, 2$, and $\varepsilon_n \leq \gamma_0$ for a subsequence $\{n_i\}_{i=1}^\infty$. However, since $(R_1, R_2) \in \mathcal{S}^*(X_1, X_2) \setminus \mathcal{V}(X_1, X_2)$ implies that $R_2 + \gamma_0 \leq H(X_2^{(1)}|X_1^{(1)}) - \gamma_0$ and $R_1 + R_2 + \gamma_0 \leq H(X_1^{(2)}X_2^{(2)}) - \gamma_0$, due to the weak law of large numbers it holds that $\Pr\{E_n^{(2)}(R_2 + \gamma_0)\} \geq 1 - \gamma_0$ for all sufficiently large odd n and $\Pr\{E_n^{(3)}(R_1 + R_2 + \gamma_0)\} \geq 1 - \gamma_0$ for all sufficiently large even n , which yields $\Pr\{E_n^{(1)}(R_1 + \gamma_0) \cup E_n^{(2)}(R_1 + \gamma_0) \cup E_n^{(3)}(R_1 + R_2 + \gamma_0)\} \geq 1 - \gamma_0$ for all sufficiently large n . Then, we can prove that $\varepsilon_n \geq 1 - \gamma_0 - 3 \cdot 2^{-n\gamma_0} \geq 1 - 2\gamma_0$ for all sufficiently large n similarly to (13) and (14) by using Lemma 2. This contradicts $\varepsilon_n \leq \gamma_0$ for all $n = n_i$, $i \geq 1$, because $\gamma_0 \in (0, \frac{1}{3})$.

5. Conclusion

In this paper the achievable rate region $\mathcal{R}_{\text{opt}}(X_1, X_2)$ in the optimistic sense for separate coding of two general sources (X_1, X_2) is discussed. We have given an inner and an outer bounds of $\mathcal{R}_{\text{opt}}(X_1, X_2)$, where the outer bound is described by using new information-theoretic quantities. We have clarified two sufficient conditions under which the inner bound coincides with the outer bounds.

We conclude this paper by pointing out that extension

of the obtained results to the case of $m (\geq 2)$ correlated general sources is straightforward.

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